

Controlling noise in the timing of intracellular events: A first–passage time approach

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Abstract

In the noisy cellular environment, gene products are subject to inherent random fluctuations in copy numbers over time. How cells ensure precision in the timing of key intracellular events, in spite of such stochasticity is an intriguing fundamental problem. We formulate event timing as a first-passage time problem, where an event is triggered when the level of a protein crosses a critical threshold for the first time. Novel analytical calculations are performed for the first-passage time distribution in stochastic models of gene expression, including models with feedback regulation. Derivation of these formulas motivates an interesting question: is there an optimal feedback strategy to regulate the synthesis of a protein to ensure that an event will occur at a precise time, while minimizing deviations or noise about the mean. Counter-intuitively, results show that for a stable long-lived protein, the optimal strategy is to express the protein at a constant rate without any feedback regulation, and any form of feedback (positive, negative or any combination of them) will always amplify noise in event timing. In contrast, a positive feedback mechanism provides the highest precision in timing for an unstable protein. These theoretical results explain recent experimental observations of single-cell lysis times in bacteriophage λ . Here, lysis of an infected bacterial cell is orchestrated by the expression and accumulation of a stable λ protein up to a threshold, and precision in timing is achieved via feedforward, rather than feedback control. Our results have broad implications for diverse cellular processes that rely on precise temporal triggering of events.

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I. INTRODUCTION

Timing of events in many cellular processes, such as cell-cycle control [1–4], cell differentiation [5, 6], sporulation [7, 8], apoptosis [9–11], development [12, 13], temporal order of gene expression [14–16], depend on regulatory proteins reaching critical threshold levels. Triggering of these events in single cells are influenced by fluctuations in protein levels that arise naturally due to noise in gene expression [17–26]. Increasing evidence shows considerable cell-to-cell variation in timing of intracellular events among isogenic cells [27–29], and it is unclear how noisy expression generates this variation. Characterization of control strategies that buffer stochasticity in event timing are critically needed to understand reliable functioning of diverse intracellular pathways that rely on precision in timing.

Mathematically, noise in the timing of events can be investigated via the first-passage time (*FPT*) framework, where an event is triggered when a stochastic process (single-cell protein level) crosses a critical threshold for the first time. There is already a rich tradition of using such first-passage time approaches to study timing of events in biological and physical sciences [30–43]. Following this tradition, exact analytical expressions for the *FPT* distribution are computed in experimentally validated and commonly used stochastic models of gene expression. These results provide novel insights into how expression parameters shape statistical fluctuations in event timing.

To investigate control mechanisms for buffering noise in timing, we consider feedback regulation in protein synthesis, where the expression rate varies arbitrarily with the protein count. Such feedback can be implemented directly through auto-regulation of gene promoter activity by its own protein [44–47] or indirectly via intermediate states [48]. It is important to point out that while the effects of such feedback loops on fluctuations in protein copy number are well studied [49–53], their impacts on stochasticity in event timing have been overlooked. We specifically formulate the problem of controlling precision in event timing as follows: what optimal form of feedback regulation ensures a given mean time to an event, while minimizing deviations or noise about the mean. It turns out that for certain minimal models of stochastic expression this optimization problem can be solved analytically, providing counter-intuitive insights. For example, in many cases negative feedback regulation is found to amplify noise in event timing, and in some cases, the optimal form of feedback is to not have any feedback at all. The robustness of these results are analyzed in the context

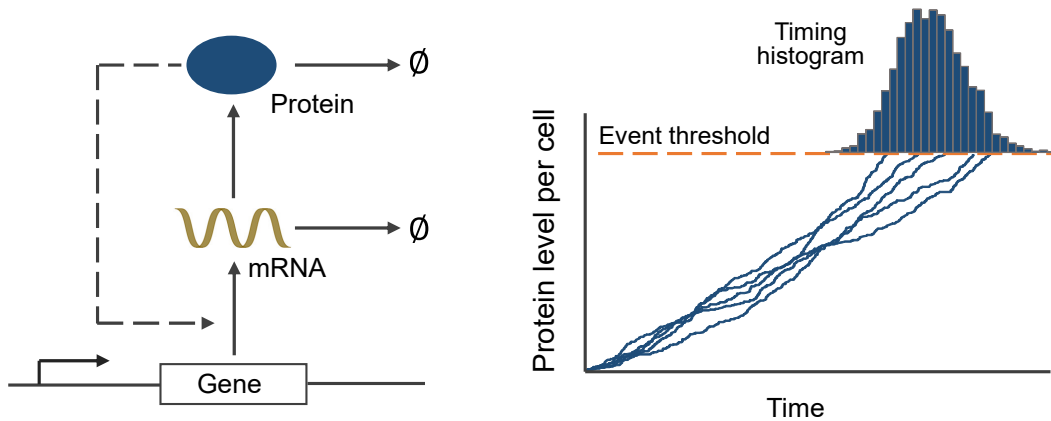


FIG. 1. **Modeling event timing as a first-passage time problem.** *Left:* Model schematic with a gene transcribing mRNAs, which is further translated into proteins. The rate of transcription is assumed to be regulated by the protein level, creating a feedback loop. *Right:* The timing of an intracellular event is formulated as the first-passage time for the protein level to reach a critical threshold. Sample trajectories for the protein level over time obtained via Monte Carlo simulations are shown, and they cross the threshold at different times due to stochasticity in gene expression. The histogram of the event timing based on 5,000 trajectories is shown on the top.

of different noise mechanisms, such as intrinsic versus extrinsic noise in gene expression [54–57]. Finally, we discuss in detail how our results explain recent experimental observations of single-cell lysis times in bacteriophage λ , where precision in timing is obtained without any feedback regulation.

II. STOCHASTIC MODEL FORMULATION

Consider a gene that is switched on at time $t = 0$ and begins to express a timekeeper protein. The intracellular event of interest is triggered once the protein reaches a critical level in the cell. We describe a minimal model of protein synthesis that incorporates two key features – expression in random bursts and feedback regulation (Fig. 1). Let $x(t) \in \{0, 1, \dots\}$ denote the level of a protein in a single cell at time t . When $x(t) = i$, the gene is transcribed at a Poisson rate k_i . Any arbitrary form of feedback can be realized by appropriately defining k_i as a function of i . For example, increasing (decreasing) k_i 's correspond to a positive (negative) feedback loop in protein production, and a fixed transcription rate implies no feedback. Assuming short-lived mRNAs, each mRNA degrades instantaneously after

producing a burst of B protein molecules [58–62]. In agreement with experimental and theoretical studies [63–65], B is assumed to follow a geometric distribution

$$\mathbb{P}(B = j) = \frac{b^j}{(b + 1)^{j+1}}, \quad b \in (0, \infty), \quad j = \{0, 1, 2, \dots\}, \quad (1)$$

where b denotes the mean protein burst size and the symbol \mathbb{P} is the notion for probability. Finally, each protein molecule degrades with a constant rate γ . The time evolution of $x(t)$ is described through the following probabilities of occurrences of burst and decay events in the next infinitesimal time $(t, t + dt]$

$$\mathbb{P}(x(t + dt) = i + B | x(t) = i) = k_i dt, \quad (2a)$$

$$\mathbb{P}(x(t + dt) = i - 1 | x(t) = i) = i\gamma dt. \quad (2b)$$

Note that in this representation of gene expression as a bursty birth-death process, the mRNA transcription rate k_i is the burst arrival rate, while the rate at which proteins are translated from an mRNA determines the mean protein burst size b . Next, we formulate event timing through the first-passage time framework.

III. COMPUTING EVENT TIMING DISTRIBUTION

The time to an event is the first-passage time for $x(t)$ to reach a threshold X starting from a zero initial condition $x(0) = 0$ (Fig. 1). It is mathematically described by the following random variable

$$FPT := \min\{t : x(t) \geq X | x(0) = 0\}, \quad (3)$$

and can be interpreted as the time taken by a random walker to first reach a defined point. Our goal is to obtain closed-form expressions for the FPT statistics in terms of underlying model parameters. Note that if the protein did not decay, then $x(t)$ accumulates over time and the FPT distribution is obtained by observing

$$\mathbb{P}(x(t) \geq X) = \mathbb{P}(FPT \leq t). \quad (4)$$

However, with protein degradation, the FPT calculation needs careful consideration so as to avoid counting multiple crossings of the threshold.

To compute the FPT we imagine the bursty birth-death process on a finite state-space $[0, 1, \dots, X]$, where the states represent the protein count (Fig. 2). All states denoting

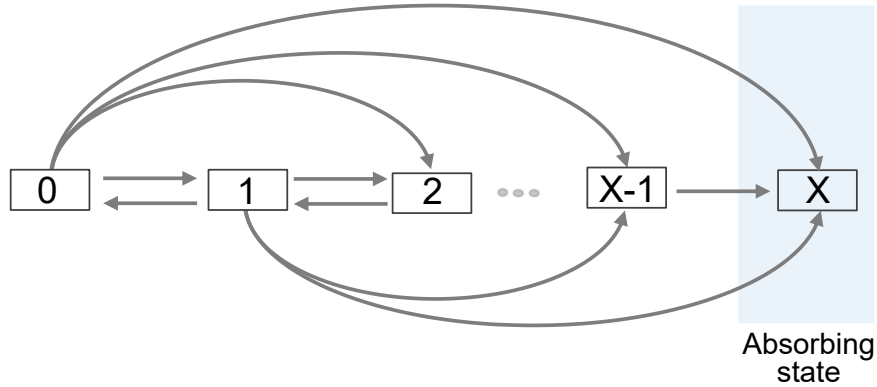


FIG. 2. **Illustration of a bursty birth-death process for computing the first-passage time:** States $[0, 1, \dots, X]$ represent the protein population counts, and arrows represent transition between states due to burst and decay events. The destination of a forward jump (a birth event) is decided by the burst size while each degradation event reduces the protein count by one. The process terminates when the protein level reaches the absorbing-state X and the first-passage time is recorded.

$x(t) \geq X$ are combined into a single absorbing state X . In this model, the probability of the protein level reaching X in the small time window $(t, t + dt)$ is the probability of being in state i at time t , and a jump of size $X - i$ or larger occurs in $(t, t + dt)$. Using the fact that for a geometrically distributed burst size B ,

$$\mathbb{P}(B \geq X - i) = \left(\frac{b}{b+1}\right)^{X-i} \quad (5)$$

and rate of burst arrival is k_i when $x(t) = i$, the probability density function (pdf) of the first-passage time is giving by

$$f_{FPT}(t) = \sum_{i=0}^{X-1} k_i \left(\frac{b}{b+1}\right)^{X-i} p_i(t) dt, \quad (6a)$$

$$p_i(t) = \mathbb{P}(x(t) = i) \quad (6b)$$

The pdf can be compactly written as a product

$$f_{FPT}(t) = \mathbf{U}\mathbf{P}(t), \quad (7a)$$

$$\mathbf{U} = \left[k_0 \left(\frac{b}{b+1}\right)^X \quad k_1 \left(\frac{b}{b+1}\right)^{X-1} \quad \dots \quad k_{X-1} \frac{b}{b+1} \right], \quad (7b)$$

$$\mathbf{P}(t) = \left[p_0(t) \quad p_1(t) \quad \dots \quad p_{X-1}(t) \right]^T, \quad (7c)$$

where \mathbf{U} is a row vector of $k_i \mathbb{P}(B \geq X - i)$ and $\mathbf{P}(t)$ is a column vector of $p_i(t)$. The time evolution of $\mathbf{P}(t)$ is given by the linear dynamical system

$$\dot{\mathbf{P}} = \mathbf{A}\mathbf{P} \quad (8)$$

derived from the Chemical Master Equations (CME) corresponding to the bursty birth-death process [66, 67]. It turns out that, in this case the matrix \mathbf{A} is a Hessenberg matrix whose i^{th} row and j^{th} column element is given by

$$a_{ij} = \begin{cases} 0, & j > i + 1 \\ (i - 1)\gamma, & j = i + 1 \\ -k_{i-1} \frac{b}{b+1} - (i - 1)\gamma, & j = i \\ k_{i-1} \frac{b^{i-j}}{(b+1)^{i-j+1}}, & j < i \end{cases} \quad (9)$$

$i, j \in \{1, \dots, X\}$. Solving (8) and using (7a) yields the following pdf for the first-passage time

$$f_{FPT}(t) = \mathbf{U}\mathbf{P}(t) = \mathbf{U} \exp(\mathbf{A}t)\mathbf{P}(0), \quad (10)$$

where $\mathbf{P}(0) = [1 \ 0 \ \dots \ 0]^T$ is vector of probabilities at $t = 0$ that follows from $x(0) = 0$. While this pdf provides complete characterization of the event timing, we are particularly interested in the lower-order statistical moments of FPT . Next, we exploit the structure of the \mathbf{A} matrix to obtain analytical formulas for the first and second order moments of the first-passage time.

IV. MOMENTS OF THE FIRST-PASSAGE TIME

From (10), the m^{th} order uncentered moment of the first-passage time is given by

$$\langle FPT^m \rangle = \mathbf{U} \left(\int_0^\infty t^m \exp(\mathbf{A}t) \right) \mathbf{P}(0). \quad (11)$$

Since the matrix \mathbf{A} is full-rank with negative eigenvalues (see section S1 of the SI), the above integral can be computed as

$$\langle FPT^m \rangle = (-1)^{m+1} m! \mathbf{U} (\mathbf{A}^{-1})^{m+1} \mathbf{P}(0). \quad (12)$$

Using the inverse of a Hessenberg matrix, we obtain the following first two moments of FPT (see section S2 of the SI)

$$\langle FPT \rangle = \frac{1}{k_0} + \frac{1}{b} \sum_{i=1}^X \frac{k_{i-1}}{k_{i-1} + (i-1)\gamma} \left(\frac{1}{k_{i-1}} + \sum_{j=i+1}^X \frac{1}{k_{j-1}} \prod_{l=i}^{j-1} \frac{k_l}{k_l + l\gamma} \frac{(b+1)l\gamma}{bk_{l-1}} \right), \quad (13a)$$

$$\langle FPT^2 \rangle = \frac{2}{b^2} \sum_{i=1}^X \frac{k_{i-1}}{k_{i-1} + (i-1)\gamma} \left(\frac{\xi_i}{k_{i-1}} + \sum_{r=i+1}^X \frac{\xi_r}{k_{r-1}} \prod_{l=i}^{r-1} \frac{k_l}{k_l + l\gamma} \frac{(b+1)l\gamma}{bk_{l-1}} \right), \quad (13b)$$

where δ_{i-1} represents the Kronecker delta which is one if $i = 1$ and zero otherwise, and

$$\xi_i = b\eta_i + \sum_{j=1}^i \eta_j, \eta_i = \frac{b}{k_0} \delta_{i-1} + \frac{k_{i-1}}{k_{i-1} + (i-1)\gamma} \left(\frac{1}{k_{i-1}} + \sum_{j=i+1}^X \frac{1}{k_{j-1}} \prod_{l=i}^{j-1} \frac{k_l}{k_l + l\gamma} \frac{(b+1)l\gamma}{bk_{l-1}} \right). \quad (14)$$

These results represent the first analytical computations of the FPT statistics for a bursty-birth death process with a random burst size and a state-dependent burst arrival rate (i.e., feedback regulation in transcription).

We investigate the complex formulas in (13) for some limiting cases. For a stable long-lived protein ($\gamma = 0$) and a constant transcription rate (no feedback; $k_i = k$), moment expressions reduce to

$$\langle FPT \rangle = \frac{1}{k} \left(\frac{X}{b} + 1 \right) \approx \frac{X}{bk}, \quad (15a)$$

$$CV_{FPT}^2 = \frac{b^2 + X + 2bX}{(b+X)^2} \approx \frac{1+2b}{X}, \quad (15b)$$

where CV_{FPT}^2 represent the noise in the first-passage time as quantified by its coefficient of variation squared. The approximate formulas in (15) are valid for a high event threshold compared to the mean protein burst size ($X/b \gg 1$). These formulas reveal important insights, such as, the noise in FPT is invariant of the transcription rate k . Moreover, $\langle FPT \rangle$ and CV_{FPT}^2 can be independently tuned — increasing the event threshold and/or reducing the burst size will lower the noise level. Once CV_{FPT}^2 is sufficient reduced, k can be modulated to adjust the mean event timing to any desired value. Next, we explore how feedback regulation of the transcription rate impacts noise in timing, for a given X and b .

V. OPTIMAL FEEDBACK STRATEGY

Having derived the FPT moments, we investigate optimal forms of transcriptional feedback that schedule an event at a given time with the highest precision. Mathematically, this

corresponds to an constraint optimization problem: find transcription rates k_0, k_1, \dots, k_{X-1} that minimize $\langle FPT^2 \rangle$ for a fixed $\langle FPT \rangle$. We first consider a stable protein whose half-life is much longer than the event timescale, and hence, degradation can be ignored ($\gamma = 0$).

A. Optimal feedback for a stable protein

When the protein of interest does not decay ($\gamma = 0$), the expressions for the FPT moments take much simpler forms

$$\begin{aligned} \langle FPT \rangle &= \frac{1}{k_0} + \frac{1}{b} \sum_{i=0}^{X-1} \frac{1}{k_i} \\ &= \frac{1}{k_0} \left(1 + \frac{1}{b} \right) + \frac{1}{k_1 b} + \frac{1}{k_2 b} \dots \end{aligned} \quad (16a)$$

$$\langle FPT^2 \rangle = \frac{2}{b^2} \left(\frac{\tau_0}{bk_0} + \sum_{i=0}^{X-1} \frac{\tau_i}{k_i} \right), \quad \tau_i := \frac{b}{k_i} + \sum_{j=i}^{X-1} \frac{1}{k_j}. \quad (16b)$$

Note that in (16a) the contribution of k_0 (transcription rate when there is no protein) is quite different from the other transcription rates $k_i, i \in \{1, 2, \dots, X-1\}$. For instance, when the event threshold is large compared to the mean burst size ($X \gg b$), then the term $1/k_0$ can be ignored and $\langle FPT \rangle \approx \sum_{i=0}^{X-1} 1/bk_i$. In contrast, if the burst size is large ($b \gg X$) then $\langle FPT \rangle \approx 1/k_0$, as a single burst event starting from zero protein molecules is sufficient for threshold crossing. Similar observation for different contributions of k_0 can be made about (16b).

It turns out that, for these simplified formulas, the problem of minimizing $\langle FPT^2 \rangle$ given $\langle FPT \rangle$ can be solved analytically using the method of Lagrange multipliers (see section S3 of the SI for details). The optimal transcription rates are given by

$$k_0 = \frac{1+b}{1+2b} \frac{2b+X}{b \langle FPT \rangle}, \quad (17a)$$

$$k_i = \frac{1+2b}{1+b} k_0 = \frac{2b+X}{b \langle FPT \rangle}, \quad \text{for } 1 \leq i \leq X-1, \quad (17b)$$

and all rates are equal to each other except for k_0 . Intuitively, the difference for k_0 comes from the fact that it contributes differently to the FPT moments as compared to other rates. Note that for a small mean burst size ($b \ll 1$), $k_0 = k_i$, whereas $k_0 = k_i/2$ for a sufficiently large b . Despite this slight deviation in k_0 , for the purposes of practical implementation,

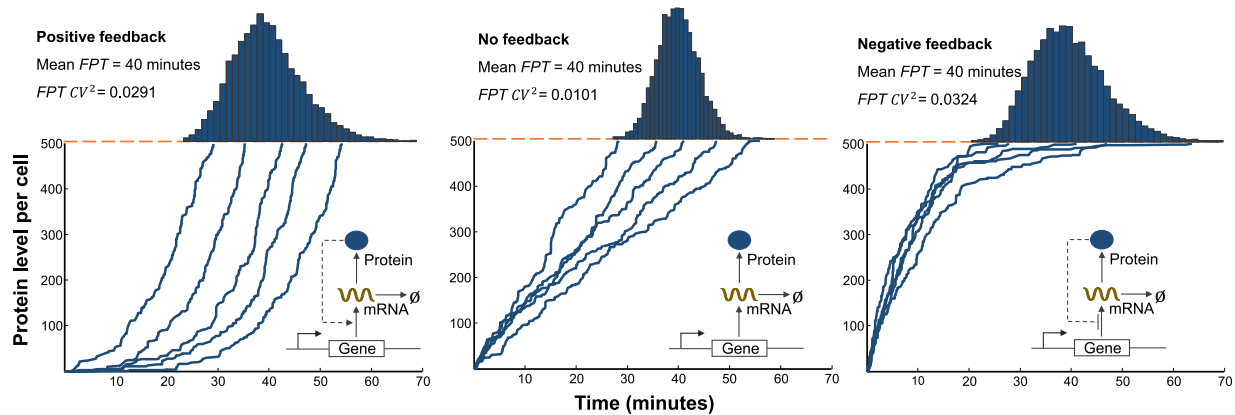


FIG. 3. **For a stable protein, no feedback provides the lowest noise in event timing for a fixed mean FPT .** Protein trajectories obtained using the Stochastic Simulation Algorithm (SSA) for a stochastic gene expression model with positive feedback (left), no feedback (middle), negative feedback (right) [68]. The threshold for event timing is assumed to be 500 protein molecules and feedback is implemented by assuming a transcription rate of the form $k_i = c_1 \pm c_2 i$, where $c_2 = 0$ (no feedback), $c_2 = 0.05 \text{ min}^{-1}$ (positive feedback), $c_2 = -0.05 \text{ min}^{-1}$ (negative feedback). For a given value of c_2 , the mean FPT is kept constant (40 minutes) by changing the parameter c_1 . The mRNA half-life is assumed to be 2.7 min, and proteins are translated from mRNAs at a rate 0.5 min^{-1} , which corresponds to a mean burst size of $b = 2$. Histograms on the top represent distribution of FPT from 10,000 Monte Carlo simulations.

the optimal feedback strategy in this case is to have a constant transcription rate (i.e., no feedback in protein expression).

We illustrate the above result via Monte Carlo simulations of stochastic gene expression models that explicitly include mRNA dynamics (Fig. 3). Here, feedback is implemented using linear transcription rates

$$k_i = c_1 + c_2 i, \quad (18)$$

where $c_2 = 0$ (no feedback), $c_2 > 0$ (positive feedback), $c_2 < 0$ (negative feedback), and $|c_2|$ is referred to as the feedback strength. As expected from theory, a no feedback strategy outperforms negative/positive feedbacks in terms of minimizing deviations in FPT around a given mean event time (Fig. 3). While similar results were obtained for implementing transcription rates using Hill functions (see section S4 of the SI), we prefer to use linearized rates as they have a fewer number of parameter and a clearer notion of feedback strength.

B. Optimal feedback for an unstable protein

Now consider the scenario where protein degradation cannot be ignored over the event timescale ($\gamma \neq 0$). Unfortunately, the expressions of the FPT moments in (13) are too convoluted for the optimization problem to be solved analytically, and the effect of different feedbacks is investigated numerically. Our strategy is as follows: choose a certain feedback strength c_2 in (18), appropriately tune c_1 using (13b) for the desired mean event timing, and explore the corresponding noise in FPT as measured by its coefficient of variation squared CV_{FPT}^2 (additional details in section S5 of the SI). Counter-intuitively, results show that for a given value of γ , a negative feedback loop in gene expression has the highest CV_{FPT}^2 , and its performance deteriorates with increasing feedback strength (Fig. 4A). In contrast, CV_{FPT}^2 first decreases with increasing strength of the positive feedback, and then increases after an optimal feedback strength is crossed (Fig. 4A). Thus, when the protein is not stable, precision in timing is attained by having a positive feedback in protein synthesis with an intermediate strength.

We next explore how the minimal achievable noise in event timing, for a fixed $\langle FPT \rangle$, varies with the protein decay rate γ . Our analysis shows that the minimum CV_{FPT}^2 obtained via positive feedback increases monotonically with γ , and $CV_{FPT}^2 \rightarrow 1$ as $\gamma \rightarrow \infty$ (Fig. 4B). Recall that the coefficient of variation of an exponentially distributed random variable is exactly equal to one. Thus, as the protein becomes more and more unstable, the timing process becomes memoryless yielding exponentially distributed first-passage times. A few interesting observations can be made from Fig. 4B: i) Higher protein burst sizes result in much larger CV_{FPT}^2 and a faster approach to $CV_{FPT}^2 = 1$ as $\gamma \rightarrow \infty$; ii) The difference in CV_{FPT}^2 for optimal feedback and no feedback is indistinguishable when the protein is stable ($\gamma = 0$) or highly unstable ($\gamma \rightarrow \infty$); and iii) For a range of intermediate protein half-lives, the optimal feedback strategy provides an order of magnitude better suppression of CV_{FPT}^2 , as compared to no feedback regulation. Taken together, these observations suggest that the timescale of protein turnover and the extent of bursty expression sets a fundamental limit to how much statistical fluctuations in FPT can be buffered.

Consistent with our above analysis, we find that a positive feedback mechanism provide the highest precision in timing in more complex stochastic expression models that explicitly include mRNA dynamics (Fig. 4C). An interesting point to note is that the protein trajec-

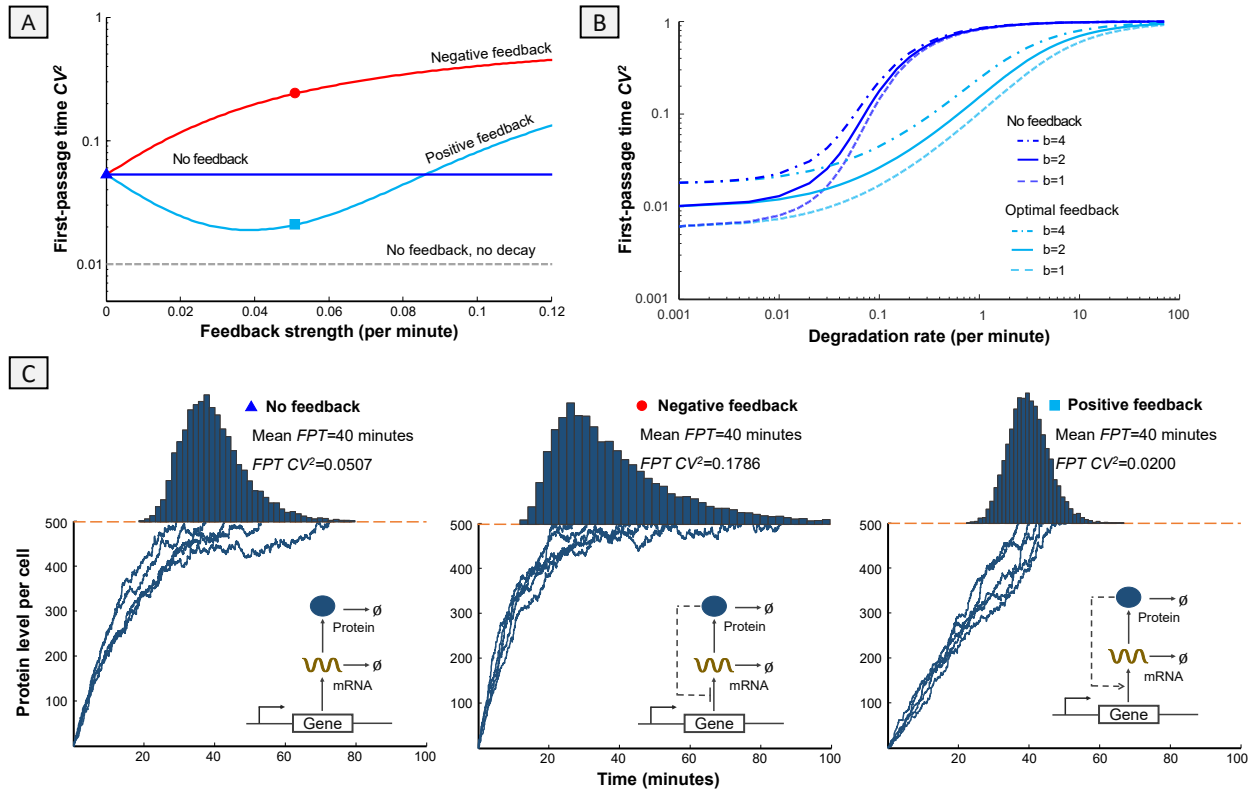


FIG. 4. **For an unstable protein, positive feedback provides the lowest noise in event timing for a fixed mean FPT .** (A) Noise in timing (CV_{FPT}^2) as a function of the feedback strength $|c_2|$ for different control strategies. The value of c_1 is changed in (18) so as to keep $\langle FPT \rangle = 40$ mins fixed. The performance of the negative feedback worsens with increasing feedback strength. In contrast, positive feedback with an optimal value of c_2 provides the highest precision in event timing. Other parameters used are $\gamma = 0.05 \text{ min}^{-1}$, $X = 500$ molecules, $b = 2$. The optimal value attained via positive feedback is much higher than the minimal value of CV_{FPT}^2 for a stable protein (dashed line). (B) The minimum CV_{FPT}^2 obtained via positive feedback increases monotonically with the protein degradation rate. For comparison purposes, CV_{FPT}^2 obtained without any feedback ($c_2 = 0$) is also plotted. (C) Monte Carlo simulations of protein trajectories under different forms of feedback control. A positive feedback outperforms the other feedback strategies in terms of minimizing statistical fluctuations in event timing. Histograms of the first-passage times obtained from 10,000 Monte Carlo runs are shown on top. For these simulations $X = 500$ molecules, $\gamma = 0.05 \text{ min}^{-1}$, mRNA half-life is assumed to be 2.7 min, proteins are translated from mRNAs at a rate 0.5 min^{-1} , and $c_2 = 0$ (no feedback), $c_2 = 0.05 \text{ min}^{-1}$ (positive feedback), $c_2 = -0.05 \text{ min}^{-1}$ (negative feedback).

tories for the optimal positive feedback case look fairly linear, and similar to the trajectories seen in the no feedback case when $\gamma = 0$ (compare rightmost plot in Fig. 4C with middle plot in Fig. 3). One way to think about this is to consider protein synthesis in the deterministic limit described by the following ordinary differential equation

$$\frac{dx(t)}{dt} = b(c_1 + c_2x) - \gamma x, \quad x(0) = 0 \quad (19)$$

with mean protein burst size b and a linear feedback form (18). If the feedback strength is chosen as

$$c_2 \approx \frac{\gamma}{b} \quad (20)$$

then the time evolution of $x(t)$ would be linear over time in (19). Indeed, our detailed stochastic analysis shows that the optimal feedback strength that minimizes CV_{FPT}^2 in the stochastic model is qualitatively similar but not identical to (20) (see section S5 of the SI).

VI. DISCUSSION

We have systematically investigated ingredients essential for precision in timing of biochemical events at the level of single cells. Our approach relies on modeling even timing as the first-passage time for a stochastically expressed protein to cross a threshold level. This framework was used to uncover optimal strategies for synthesizing the protein that ensures a given mean time to event triggering (threshold crossing), with minimal fluctuations around the mean. The main contributions and insights can be summarized as:

1. Novel analytical calculations for the first-passage time in stochastic models of gene expression, with and without feedback regulation.
2. If the protein half-life is much longer than the timescale of the event, the highest precision in event timing is attained by having no feedback, i.e., express the protein at a constant rate (Fig. 3).
3. In the absence of feedback, the noise in event timing is given by (15) and determined by the molecular threshold (X) and the protein burst size (b). Once X and b are chosen for a tolerable noise level, the mean time to the event can be adjusted independently through the transcription rate k .

4. If the protein half-life is comparable or shorter than the timescale of the event, then positive feedback provides the lowest noise in event timing (Fig. 4A). Moreover, negative feedback always amplifies noise around the mean time.
5. The minimum achievable noise in timing increases with the protein decay rate γ and approaches $CV_{FPT}^2 = 1$ as $\gamma \rightarrow \infty$ (Fig. 4B).

How robust are these findings to alternative noise sources and key modeling assumptions? For example, the model only considers noise from low-copy number fluctuation in gene product levels, and ignores any form of “extrinsic noise” that arises from cell-to-cell differences in gene expression machinery [55, 69]. To incorporate such extrinsic noise, we alter the transcription rate to $k_i Z$, where Z is drawn from an a priori probability distribution at the start of gene expression ($t = 0$), and remains fixed till the threshold is reached. Interestingly, the optimal feedback derived in (17) does not change even after adding extrinsic noise to the transcription rate or the protein burst size (see section S6 of the SI). Another important model feature is geometrically distributed protein burst sizes, which follows from the assumption of exponentially distributed mRNA lifetimes. We have also explored the scenario of perfect memory in the mRNA degradation process, which results in a mRNA lifetime distribution given by the delta function. In this case, the protein burst size is Poisson and the optimal feedback strategy is fairly close to having no feedback for a stable protein (see section S7 of the SI). Next, we discuss the biological implications of our findings in the context of phage λ 's lysis times, i.e., the time taken by the virus to destroy infected bacterial cells.

A. Connecting theoretical insights to λ lysis times

Phage λ has recently emerged as a simple model system for studying event timing at the level of single cells [27, 28]. After infecting *E. coli*, λ expresses a protein, holin, which accumulates in the inner membrane. When holin reaches a critical threshold concentration, it undergoes a structural transformation, forming holes in the membrane [70]. These holes allow other lysis proteins (endolysin and spanin) to access and rupture the cell wall [71]. Subsequently the cell lysis and phage progeny are released into the surrounding medium. Since hole formation and cell rupture are nearly simultaneous, lysis timing depends on de

novo expression and accumulation of holin in the cell membrane up to a critical threshold [70–72]. Data reveals precision in the timing of lysis – individual cells infected by a single virus lyse on average at 65 mins, with a standard deviation of 3.5 mins, implying a coefficient of variation of $\approx 5\%$. Such precision is expected given the existence of an optimal lysis time [73–77]. Intuitively, if λ lysis is early then there are no viral progeny. In contrast, if λ lysis is late then the infected cell could die before lysis is effected, trapping the virus with it.

The threshold for lysis is reported to be a few thousand holin molecules [78]. Moreover, the holin mean burst size (average number of holins produced in a single mRNA lifetime) is estimated as $b \approx 1 - 3$ [78]. Based on our *FPT* moment calculations in (15), such a small protein burst size relative to the event threshold will yield a tight distribution of lysis times. Interestingly, (15) provides insights for engineering mutant λ that lyse, on average, at the same time as the wild type, but with much higher noise. This could be done by lowering the threshold for lysis through mutations in the holin amino acid sequence [28], and also reducing the holin mRNA transcription or translation rate so as to keep the same mean lysis time. It is important to point out that since holin proteins are long-lived and do not degrade over relevant timescales [79], λ 's lysis system with no known feedback in holin expression provides better suppression of lysis-time fluctuation compared to any feedback regulated system.

B. Additional mechanism for noise buffering

The surprising ineffectiveness of feedback control motivates the need for other mechanisms to buffer noise in event timing. Intriguingly, λ uses feedforward control to regulate the timing of lysis. Feedforward control is implemented through two proteins with opposing functions: holin and antiholin [80–82]. In the wild-type virus both proteins are expressed in a 2:1 ratio (for every two holins there is one antiholin) from the same mRNA through a dual start motif. Antiholin binds to holin and prevents holin from participating in hole formation, creating an incoherent feedforward circuit (Fig. 5). Synthesis of antiholin leads to a lower burst size for active holin molecules, and increases the threshold for the total number of holins needed for lysis – both factors functioning to lower the noise in event timing. Consistent with this prediction, variants of λ lacking antiholin are experimentally observed to exhibit much higher intercellular variation in lysis times as compared to the

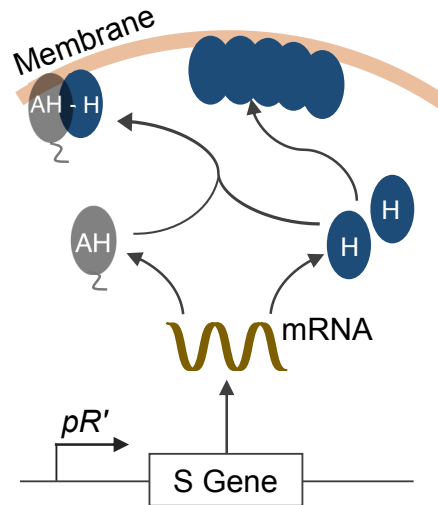


FIG. 5. **An incoherent feedforward circuit in λ 's lytic pathway:** Bacteriophage λ lysis the infected host cell by expressing a membrane protein, holin (H). The protein slowly accumulates on the cell membrane over time and forms holes when a critical concentration threshold is reached. The mRNA encoding holin also expresses antiholin (AH), which binds to holin and prevents it from participating in hole formation creating a feedforward circuit.

wild-type virus [28, 83]. In summary, λ encodes a multitude of regulatory mechanisms (low holin burst size; no feedback regulation; feedforward control) to ensure that single infected cells lyse at an optimal time, in spite of the inherently stochastic expression of lysis proteins. These results illustrate the utility of the first-passage time framework for characterizing noise in the timing of intracellular events. Finally, analytical results and insights obtained here have broader implications for timing phenomenon in chemical kinetics, ecological modeling and statistical physics.

SUPPLEMENTARY INFORMATION

S1. ON SOME PROPERTIES OF THE MATRIX \mathbf{A}

In this section, we discuss some properties of the matrix \mathbf{A} given in equation (9) of the main text.

S1-a. \mathbf{A} is a Hurwitz matrix

The matrix \mathbf{A} is given by

$$\mathbf{A} = \begin{bmatrix} -\frac{bk_0}{b+1} & \gamma & \cdots & 0 \\ \frac{bk_0}{(b+1)^2} & -\left(\frac{bk_1}{b+1} + \gamma\right) & \cdots & 0 \\ \vdots & \vdots & \vdots & \vdots \\ \frac{b^{X-2}k_0}{(b+1)^{X-1}} & \frac{b^{X-3}k_1}{(b+1)^{X-2}} & \cdots & (X-1)\gamma \\ \frac{b^{X-1}k_0}{(b+1)^X} & \frac{b^{X-2}k_1}{(b+1)^{X-1}} & \cdots & -\left(\frac{bk_{X-1}}{b+1} + (X-1)\gamma\right) \end{bmatrix}. \quad (\text{S1.1})$$

In order to prove that \mathbf{A} is a Hurwitz matrix, we prove that the following two conditions hold true [84, pp. 48–49]:

1. The diagonal elements $a_{ii} < 0$ for $i = 1, 2, \dots, X$,

$$2. \max_{1 \leq j \leq X} \sum_{\substack{i=1 \\ j \neq i}}^X \left| \frac{a_{ij}}{a_{jj}} \right| < 1.$$

Condition 1 here is satisfied here as $a_{ii} = -\frac{bk_{i-1}}{b+1} - (i-1)\gamma < 0$. The left hand side of Condition 2 for any column $j = 1, 2, \dots, X$ is

$$\sum_{\substack{i=1 \\ j \neq i}}^X \left| \frac{a_{ij}}{a_{jj}} \right| = \frac{(j-1)\gamma}{\frac{bk_{j-1}}{b+1} + (j-1)\gamma} + \frac{bk_{j-1}}{b+1} \sum_{i=j+1}^X \frac{\left(\frac{b}{b+1}\right)^{i-j}}{\frac{bk_{j-1}}{b+1} + (j-1)\gamma} \quad (\text{S1.2})$$

$$= \frac{(j-1)\gamma}{\frac{bk_{j-1}}{b+1} + (j-1)\gamma} + \frac{\frac{bk_{j-1}}{b+1} \left(1 - \left(\frac{b}{b+1}\right)^{X-j}\right)}{\frac{bk_{j-1}}{b+1} + (j-1)\gamma} \quad (\text{S1.3})$$

$$= \frac{\frac{bk_{j-1}}{b+1} \left(1 - \left(\frac{b}{b+1}\right)^{X-j}\right) + (j-1)\gamma}{\frac{bk_{j-1}}{b+1} + (j-1)\gamma} < 1. \quad (\text{S1.4})$$

Thus, the matrix \mathbf{A} is Hurwitz, i.e., the eigenvalues of \mathbf{A} have negative real part.

S1-b. \mathbf{A} is an invertible matrix

We will find the inverse of the matrix \mathbf{A} thus proving that it is invertible.

Let us use \mathbf{A}_0 to denote the matrix \mathbf{A} when $\gamma = 0$. The lower triangular matrix \mathbf{A}_0 is given by

$$\mathbf{A} = \begin{bmatrix} -\frac{bk_0}{b+1} & 0 & \cdots & 0 \\ \frac{bk_0}{(b+1)^2} & -\frac{bk_1}{b+1} & \cdots & 0 \\ \vdots & \vdots & \vdots & \vdots \\ \frac{b^{X-2}k_0}{(b+1)^{X-1}} & \frac{b^{X-3}k_1}{(b+1)^{X-2}} & \cdots & 0 \\ \frac{b^{X-1}k_0}{(b+1)^X} & \frac{b^{X-2}k_1}{(b+1)^{X-1}} & \cdots & -\frac{bk_{X-1}}{b+1} \end{bmatrix}. \quad (\text{S1.5})$$

We claim that the inverse of \mathbf{A}_0 is given by the following matrix

$$\mathbf{A}_0^{-1} = -\frac{1}{b} \begin{bmatrix} \frac{b+1}{k_0} & 0 & \cdots & 0 & 0 \\ \frac{1}{k_1} & \frac{b+1}{k_1} & \cdots & 0 & 0 \\ \frac{1}{k_2} & \frac{1}{k_2} & \cdots & 0 & 0 \\ \vdots & \vdots & \vdots & \vdots & \vdots \\ \frac{1}{k_{X-2}} & \frac{1}{k_{X-2}} & \cdots & \frac{b+1}{k_{X-2}} & 0 \\ \frac{1}{k_{X-1}} & \frac{1}{k_{X-1}} & \cdots & \frac{1}{k_{X-1}} & \frac{b+1}{k_{X-1}} \end{bmatrix}. \quad (\text{S1.6})$$

This claim can be quickly verified by multiplying the matrices which results in identity matrix. Next, to determine \mathbf{A}^{-1} , we observe that when $\gamma \neq 0$, the matrix \mathbf{A} can be written as

$$\mathbf{A} = \mathbf{A}_0 + \mathbf{A}_e, \quad (\text{S1.7})$$

where \mathbf{A}_e is given by

$$\mathbf{A}_e = \begin{bmatrix} 0 & \gamma & \cdots & 0 & 0 \\ 0 & -\gamma & \cdots & 0 & 0 \\ 0 & 0 & \ddots & 0 & 0 \\ \vdots & \vdots & \vdots & \vdots & \vdots \\ 0 & 0 & \cdots & -(X-2)\gamma & (X-1)\gamma \\ 0 & 0 & \cdots & 0 & -(X-1)\gamma \end{bmatrix}. \quad (\text{S1.8})$$

Therefore the inverse of the matrix \mathbf{A} can be written as

$$\mathbf{A}^{-1} = (\mathbf{A}_0 + \mathbf{A}_e)^{-1} = (\mathbf{I} + \mathbf{A}_0^{-1}\mathbf{A}_e)^{-1} \mathbf{A}_0^{-1}. \quad (\text{S1.9})$$

Note that \mathbf{A}_e is a bidiagonal matrix with the i^{th} diagonal element as $-(i-1)\gamma$, $1 \leq i \leq X$ while the j^{th} super-diagonal elements as $j\gamma$, $1 \leq j \leq X-1$. As we have already determined the expression of \mathbf{A}_0^{-1} in (S1.6), the expression of $\mathbf{A}_0^{-1}\mathbf{A}_e$ can be determined as

$$-\frac{1}{b} \begin{bmatrix} \frac{b+1}{k_0} & 0 & \cdots & 0 & 0 \\ \frac{1}{k_1} & \frac{b+1}{k_1} & \cdots & 0 & 0 \\ \frac{1}{k_2} & \frac{1}{k_2} & \cdots & 0 & 0 \\ \vdots & \vdots & \vdots & \vdots & \vdots \\ \frac{1}{k_{X-2}} & \frac{1}{k_{X-2}} & \cdots & \frac{b+1}{k_{X-2}} & 0 \\ \frac{1}{k_{X-1}} & \frac{1}{k_{X-1}} & \cdots & \frac{1}{k_{X-1}} & \frac{b+1}{k_{X-1}} \end{bmatrix} \begin{bmatrix} 0 & \gamma & \cdots & 0 & 0 \\ 0 & -\gamma & \cdots & 0 & 0 \\ 0 & 0 & \ddots & 0 & 0 \\ \vdots & \vdots & \vdots & \vdots & \vdots \\ 0 & 0 & \cdots & -(X-2)\gamma & (X-1)\gamma \\ 0 & 0 & \cdots & 0 & -(X-1)\gamma \end{bmatrix} \quad (\text{S1.10})$$

$$= \begin{bmatrix} 0 & -\frac{(b+1)\gamma}{bk_0} & \cdots & 0 & 0 \\ 0 & \frac{\gamma}{k_1} & \cdots & 0 & 0 \\ 0 & 0 & \ddots & 0 & 0 \\ \vdots & \vdots & \vdots & \vdots & \vdots \\ 0 & 0 & \cdots & \frac{(X-2)\gamma}{k_{X-2}} & -\frac{(X-1)(b+1)\gamma}{bk_{X-2}} \\ 0 & 0 & \cdots & 0 & \frac{(X-1)\gamma}{k_{X-1}} \end{bmatrix}. \quad (\text{S1.11})$$

Thus, the matrix $\mathbf{I} + \mathbf{A}_0^{-1}\mathbf{A}_e$ is a bidiagonal matrix with its diagonal elements $1 + \frac{(i-1)\gamma}{k_{i-1}} = \frac{k_{i-1} + (i-1)\gamma}{k_{i-1}}$ for $i = 1, 2, \dots, X$. The super diagonal elements are given by $-\frac{j(b+1)\gamma}{bk_{j-1}}$ for $j = 1, 2, \dots, X-1$. Using the result for inverse of a bidiagonal matrix derived in [85], we can write the (i, j) element of inverse of $\mathbf{E} := \mathbf{I} + \mathbf{A}_0^{-1}\mathbf{A}_e$ as follows

$$e'_{i,j} = \begin{cases} 0 & \text{if } i > j, \\ \frac{k_{i-1}}{k_{i-1} + (i-1)\gamma}, & \text{if } i = j, \\ \frac{k_{i-1}}{k_{i-1} + (i-1)\gamma} \prod_{l=i}^{j-1} \left(\frac{k_l}{k_l + l\gamma} \frac{l(b+1)\gamma}{bk_{l-1}} \right), & \text{if } i < j. \end{cases} \quad (\text{S1.12})$$

Alternatively, in matrix form

$$\mathbf{E}^{-1} = \begin{bmatrix} 1 & \frac{k_1}{k_1 + \gamma} \frac{(b+1)\gamma}{bk_0} & \cdots & \prod_{l=1}^{X-1} \left(\frac{k_l}{k_l + l\gamma} \frac{l(b+1)\gamma}{bk_{l-1}} \right) \\ 0 & \frac{k_1}{k_1 + \gamma} & \cdots & \frac{k_1}{k_1 + \gamma} \prod_{l=2}^{X-1} \left(\frac{k_l}{k_l + l\gamma} \frac{l(b+1)\gamma}{bk_{l-1}} \right) \\ \vdots & \vdots & \vdots & \vdots \\ 0 & 0 & \cdots & \frac{k_{X-1}}{k_{X-1} + (X-1)\gamma} \end{bmatrix}. \quad (\text{S1.13})$$

We can compute \mathbf{A}^{-1} by calculating $\mathbf{E}^{-1}\mathbf{A}_0^{-1}$. Here, we do not give explicit form of \mathbf{A}^{-1} as it is not required for calculations in this document. Nevertheless, it validates that \mathbf{A} is an invertible matrix.

S2. EXPRESSION OF m^{th} MOMENT OF FIRST-PASSAGE TIME

In this section, we make use of the properties discussed in the previous section to determine the moments of the first-passage time. As discussed in the main text (equation (12)), the distribution of first-passage time (*FPT*) is given by the following

$$f_{FPT}(t) = \mathbf{U} \exp(\mathbf{A}t) \mathbf{P}(0). \quad (\text{S2.1})$$

Therefore, a moment of m^{th} order can be calculated as

$$\langle FPT^m \rangle = \int_0^\infty t^m \mathbf{U} \exp(\mathbf{A}t) \mathbf{P}(0) dt \quad (\text{S2.2})$$

$$= \mathbf{U} \left(\int_0^\infty t^m \exp(\mathbf{A}t) dt \right) \mathbf{P}(0). \quad (\text{S2.3})$$

Let us consider $\int_0^\infty t^m \exp(\mathbf{A}t) dt$. Integrating by parts

$$\int_0^\infty t^m \exp(\mathbf{A}t) dt = t^m \mathbf{A}^{-1} \exp(\mathbf{A}t) \Big|_0^\infty - m \mathbf{A}^{-1} \int_0^\infty t^{m-1} \exp(\mathbf{A}t) dt. \quad (\text{S2.4})$$

The first part of on the right hand side goes to zero as long as $\exp(\mathbf{A}t)$ goes to zero faster than the polynomial term t^m which is true since the matrix \mathbf{A} is a Hurwitz matrix, i.e., the eigenvalues of \mathbf{A} have negative real parts. Using a_m as a notion to represent $\int_0^\infty t^m \exp(\mathbf{A}t) dt$, we can write the following recursive relationship

$$a_m = -m \mathbf{A}^{-1} a_{m-1}. \quad (\text{S2.5})$$

Thus, $a_m = (-1)^m m! (\mathbf{A}^{-1})^m a_0$. Further,

$$a_0 = \int_0^\infty \exp(\mathbf{A}t) dt = \mathbf{A}^{-1} \exp(\mathbf{A}t) \Big|_0^\infty = -\mathbf{A}^{-1}. \quad (\text{S2.6})$$

Therefore $a_m = (-1)^{m+1} m! (\mathbf{A}^{-1})^{m+1}$. Substituting this in (S2.3) gives the following for a general m^{th} moment of FPT

$$\langle FPT^m \rangle = (-1)^m m! \mathbf{U} (\mathbf{A}^{-1})^{m+1} \mathbf{P}(0). \quad (\text{S2.7})$$

S2-a. Calculation of $\mathbf{U} \mathbf{A}^{-1}$

As we saw in equation (S2.7), calculation of the moments will have a term of the form $\mathbf{U} \mathbf{A}^{-1}$. Here we provide the calculation of this term.

Consider two matrices \mathbf{G} and \mathbf{H} such that $\mathbf{A}_e = \mathbf{G}\mathbf{H}$ where \mathbf{G} is a $X \times X - 1$ matrix

$$\mathbf{G} = \gamma \begin{bmatrix} -1 & 0 & 0 & \cdots & 0 & 0 \\ 1 & -2 & 0 & \cdots & 0 & 0 \\ 0 & 2 & -3 & \cdots & 0 & 0 \\ 0 & 0 & 3 & \cdots & 0 & 0 \\ \vdots & \vdots & \vdots & \vdots & \vdots & \\ 0 & 0 & 0 & \cdots & X - 2 & -(X - 1) \\ 0 & 0 & 0 & \cdots & 0 & X - 1 \end{bmatrix}, \quad (\text{S2.8})$$

while \mathbf{H} is a $X - 1 \times X$ matrix

$$\mathbf{H} = \begin{bmatrix} 0 & -1 & 0 & \cdots & 0 & 0 \\ 0 & 0 & -1 & \cdots & 0 & 0 \\ \vdots & \vdots & \vdots & \vdots & \vdots & \\ 0 & 0 & 0 & \cdots & -1 & 0 \\ 0 & 0 & 0 & \cdots & 0 & -1 \end{bmatrix}. \quad (\text{S2.9})$$

Using the matrix inversion lemma, \mathbf{A}^{-1} can be written as

$$\mathbf{A}^{-1} = (\mathbf{A}_0 + \mathbf{G}\mathbf{H})^{-1}, \quad (\text{S2.10})$$

$$= \mathbf{A}_0^{-1} - \mathbf{A}_0^{-1}\mathbf{G}(\mathbf{I} + \mathbf{H}\mathbf{A}_0^{-1}\mathbf{G})^{-1}\mathbf{H}\mathbf{A}_0^{-1}. \quad (\text{S2.11})$$

Let us look at the expression $\mathbf{U}\mathbf{A}_0^{-1}\mathbf{G}$.

$$\mathbf{U}\mathbf{A}_0^{-1}\mathbf{G} = - \begin{bmatrix} 1 & 1 & \cdots & 1 \end{bmatrix} \begin{bmatrix} -1 & 0 & \cdots & 0 \\ 1 & -2 & \cdots & 0 \\ 0 & 2 & \cdots & 0 \\ 0 & 0 & \cdots & 0 \\ \vdots & \vdots & \vdots & \vdots \\ 0 & 0 & \cdots & -(X - 1) \\ 0 & 0 & \cdots & X - 1 \end{bmatrix}, \quad (\text{S2.12})$$

$$= - \begin{bmatrix} 0 & 0 & \cdots & 0 \end{bmatrix}. \quad (\text{S2.13})$$

Therefore, we can conclude that $\mathbf{U}\mathbf{A}^{-1}$ is in fact equal to $\mathbf{U}\mathbf{A}_0^{-1}$ which could be calculated by multiplying \mathbf{U} and \mathbf{A}_0^{-1} .

$$\mathbf{U}\mathbf{A}_0^{-1} = -\frac{1}{b} \begin{bmatrix} k_0 \left(\frac{b}{b+1}\right)^X \\ k_1 \left(\frac{b}{b+1}\right)^{X-1} \\ k_2 \left(\frac{b}{b+1}\right)^{X-2} \\ \vdots \\ k_{X-2} \left(\frac{b}{b+1}\right)^2 \\ k_{X-1} \left(\frac{b}{b+1}\right) \end{bmatrix}^T \begin{bmatrix} \frac{b+1}{k_0} & 0 & \dots & 0 & 0 \\ \frac{1}{k_1} & \frac{b+1}{k_1} & \dots & 0 & 0 \\ \frac{1}{k_2} & \frac{1}{k_2} & \dots & 0 & 0 \\ \vdots & \vdots & \vdots & \vdots & \vdots \\ \frac{1}{k_{X-2}} & \frac{1}{k_{X-2}} & \dots & \frac{b+1}{k_{X-2}} & 0 \\ \frac{1}{k_{X-1}} & \frac{1}{k_{X-1}} & \dots & \frac{1}{k_{X-1}} & \frac{b+1}{k_{X-1}} \end{bmatrix} \quad (\text{S2.14})$$

$$= -\frac{1}{b} \begin{bmatrix} \frac{b^X}{(b+1)^{X-1}} + \sum_{l=1}^{X-1} \left(\frac{b}{b+1}\right)^l \\ \frac{b^{X-1}}{(b+1)^{X-2}} + \sum_{l=1}^{X-2} \left(\frac{b}{b+1}\right)^l \\ \frac{b^{X-2}}{(b+1)^{X-3}} + \sum_{l=1}^{X-3} \left(\frac{b}{b+1}\right)^l \\ \vdots \\ \frac{b^2}{b+1} + \frac{b}{b+1} \\ b \end{bmatrix}^T \quad (\text{S2.15})$$

$$= -\frac{1}{b} [b \ b \ b \ \dots \ b \ b] = -[1 \ 1 \ 1 \ \dots \ 1 \ 1]. \quad (\text{S2.16})$$

S2-b. Calculation of the first-two first-passage time moments

Mean FPT The mean FPT's expression can be written as

$$\langle FPT \rangle = \mathbf{U}\mathbf{A}^{-2}\mathbf{P}(0) = \mathbf{U}\mathbf{A}_0^{-1}\mathbf{E}^{-1}\mathbf{A}_0^{-1}\mathbf{P}(0). \quad (\text{S2.17})$$

The expression of $\mathbf{A}_0^{-1}\mathbf{P}(0)$ is just the first column of \mathbf{A}_0 . Therefore

$$\mathbf{E}^{-1}\mathbf{A}_0^{-1}\mathbf{P}(0) = -\frac{1}{b} \begin{bmatrix} 1 & \frac{k_1}{k_1+\gamma} \frac{(b+1)\gamma}{bk_0} & \dots & \prod_{l=1}^{X-1} \left(\frac{k_l}{k_l+l\gamma} \frac{l(b+1)\gamma}{bk_{l-1}} \right) \\ 0 & \frac{k_1}{k_1+\gamma} & \dots & \frac{k_1}{k_1+\gamma} \prod_{l=2}^{X-1} \left(\frac{k_l}{k_l+l\gamma} \frac{l(b+1)\gamma}{bk_{l-1}} \right) \\ \vdots & \vdots & \vdots & \vdots \\ 0 & 0 & \dots & \frac{k_{X-1}}{k_{X-1}+(X-1)\gamma} \end{bmatrix} \begin{bmatrix} \frac{b+1}{k_0} \\ \frac{1}{k_1} \\ \vdots \\ \frac{1}{k_{X-1}} \end{bmatrix} \quad (\text{S2.18})$$

$$= -\frac{1}{b} \begin{bmatrix} \frac{b+1}{k_0} + \frac{1}{k_1} \frac{k_1}{k_1+\gamma} \frac{(b+1)\gamma}{bk_0} + \dots + \frac{1}{k_{X-1}} \prod_{l=1}^{X-1} \left(\frac{k_l}{k_l+l\gamma} \frac{l(b+1)\gamma}{bk_{l-1}} \right) \\ \frac{1}{k_1} \frac{k_1}{k_1+\gamma} + \frac{1}{k_2} \frac{k_1}{k_1+\gamma} \frac{k_2}{k_2+2\gamma} \frac{2(b+1)\gamma}{bk_1} + \dots + \frac{1}{k_{X-1}} \frac{k_1}{k_1+\gamma} \prod_{l=2}^{X-1} \left(\frac{k_l}{k_l+l\gamma} \frac{l(b+1)\gamma}{bk_{l-1}} \right) \\ \vdots \\ \frac{1}{k_{X-1}} \frac{k_{X-1}}{k_{X-1}+(X-1)\gamma} \end{bmatrix}. \quad (\text{S2.19})$$

Since the vector $\mathbf{U}\mathbf{A}_0^{-1} = -\begin{bmatrix} 1 & 1 & 1 & \dots & 1 & 1 \end{bmatrix}$, $\mathbf{U}\mathbf{A}_0^{-1}\mathbf{E}^{-1}\mathbf{A}_0^{-1}\mathbf{P}(0)$ is essentially negative sum of the elements of $\mathbf{E}^{-1}\mathbf{A}_0^{-1}\mathbf{P}(0)$. Therefore we have the expression of mean FPT is given by

$$\langle FPT \rangle = \frac{1}{k_0} + \frac{1}{b} \sum_{i=1}^X \left(\frac{1}{k_{i-1}} \frac{k_{i-1}}{k_{i-1} + (i-1)\gamma} + \sum_{j=i+1}^X \frac{1}{k_{j-1}} \left(\frac{k_{i-1}}{k_{i-1} + (i-1)\gamma} \prod_{l=i}^{j-1} \left(\frac{k_l}{k_l + l\gamma} \frac{l(b+1)\gamma}{bk_{l-1}} \right) \right) \right). \quad (\text{S2.20})$$

a. *Second order moment* The second order moment is given by

$$\langle FPT^2 \rangle = -2\mathbf{U}\mathbf{A}^{-3}\mathbf{P}(0) = -2\mathbf{U}\mathbf{A}_0^{-1}\mathbf{E}^{-1}\mathbf{A}_0^{-1}\mathbf{E}^{-1}\mathbf{A}_0^{-1}\mathbf{P}(0). \quad (\text{S2.21})$$

Let us use the notation η_i defined as

$$\eta_i := \frac{k_{i-1}}{k_{i-1} + (i-1)\gamma} \left(\frac{1}{k_{i-1}} + \sum_{j=i+1}^X \frac{1}{k_{j-1}} \prod_{l=i}^{j-1} \frac{k_l}{k_l + l\gamma} \frac{l(b+1)\gamma}{bk_{l-1}} \right) + \frac{b}{k_0} \delta_{i-1}. \quad (\text{S2.22})$$

Using (S2.19), we can write

$$\mathbf{E}^{-1}\mathbf{A}_0^{-1}\mathbf{P}(0) = -\frac{1}{b} \begin{bmatrix} \eta_1 \\ \eta_1 \\ \vdots \\ \eta_{X-2} \\ \eta_{X-1} \end{bmatrix}. \quad (\text{S2.23})$$

Therefore

$$\mathbf{A}_0^{-1}\mathbf{E}^{-1}\mathbf{A}_0^{-1}\mathbf{P}(0) = \left(\frac{1}{b}\right)^2 \begin{bmatrix} \frac{b+1}{k_0} & 0 & \dots & 0 & 0 \\ \frac{1}{k_1} & \frac{b+1}{k_1} & \dots & 0 & 0 \\ \vdots & \vdots & \ddots & \vdots & \vdots \\ \frac{1}{k_{X-2}} & \frac{1}{k_{X-2}} & \dots & \frac{b+1}{k_{X-2}} & 0 \\ \frac{1}{k_{X-1}} & \frac{1}{k_{X-1}} & \dots & \frac{1}{k_{X-1}} & \frac{b+1}{k_{X-1}} \end{bmatrix} \begin{bmatrix} \eta_1 \\ \eta_2 \\ \vdots \\ \eta_{X-1} \\ \eta_X \end{bmatrix} \quad (\text{S2.24})$$

$$= \left(\frac{1}{b}\right)^2 \begin{bmatrix} \frac{(b+1)\eta_1}{k_0} \\ \frac{\eta_1 + (b+1)\eta_2}{k_1} \\ \vdots \\ \frac{\sum_{i=1}^{X-2} \eta_i + (b+1)\eta_{X-1}}{k_{X-2}} \\ \frac{\sum_{i=1}^{X-1} \eta_i + (b+1)\eta_X}{k_{X-1}} \end{bmatrix}. \quad (\text{S2.25})$$

Using the notion $\xi_i = \sum_{j=1}^{i-1} \eta_j + (b+1)\eta_i$, we can write $\mathbf{E}^{-1}\mathbf{A}_0^{-1}\mathbf{E}^{-1}\mathbf{A}_0^{-1}\mathbf{P}(0)$ as

$$\left(\frac{1}{b}\right)^2 \begin{bmatrix} 1 & \frac{k_1}{k_1+\gamma} \frac{(b+1)\gamma}{bk_0} & \cdots & \prod_{l=1}^{X-1} \left(\frac{k_l}{k_l+l\gamma} \frac{l(b+1)\gamma}{bk_{l-1}} \right) \\ 0 & \frac{k_1}{k_1+\gamma} & \cdots & \frac{k_1}{k_1+\gamma} \prod_{l=2}^{X-1} \left(\frac{k_l}{k_l+l\gamma} \frac{l(b+1)\gamma}{bk_{l-1}} \right) \\ \vdots & \vdots & \vdots & \vdots \\ 0 & 0 & \cdots & \frac{k_{X-1}}{k_{X-1}+(X-1)\gamma} \end{bmatrix} \begin{bmatrix} \xi_1 \\ k_0 \\ \xi_2 \\ k_1 \\ \vdots \\ \xi_X \\ k_{X-1} \end{bmatrix} \quad (\text{S2.26})$$

$$= \left(\frac{1}{b}\right)^2 \begin{bmatrix} \frac{\xi_1}{k_0} + \frac{\xi_2}{k_1} \frac{k_1}{k_1+\gamma} \frac{(b+1)\gamma}{bk_0} + \cdots + \frac{\xi_X}{k_{X-1}} \prod_{l=1}^{X-1} \left(\frac{k_l}{k_l+l\gamma} \frac{l(b+1)\gamma}{bk_{l-1}} \right) \\ \frac{\xi_2}{k_1} \frac{k_1}{k_1+\gamma} + \cdots + \frac{\xi_X}{k_{X-1}} \frac{k_1}{k_1+\gamma} \prod_{l=2}^{X-1} \left(\frac{k_l}{k_l+l\gamma} \frac{l(b+1)\gamma}{bk_{l-1}} \right) \\ \vdots \\ \frac{\xi_X}{k_{X-1}} \frac{k_{X-1}}{k_{X-1}+(X-1)\gamma} \end{bmatrix}. \quad (\text{S2.27})$$

As $\mathbf{U}\mathbf{A}_0^{-1}$ is merely negative summation of the elements of the column vector it pre-multiplies to, the second order moment of FPT can be given by following explicit formula

$$\langle FPT^2 \rangle = 2 \left(\frac{1}{b}\right)^2 \sum_{i=1}^X \frac{k_{i-1}}{k_{i-1} + (i-1)\gamma} \left(\frac{\xi_i}{k_{i-1}} + \sum_{r=i+1}^X \frac{\xi_r}{k_{r-1}} \left(\prod_{l=i}^{r-1} \left(\frac{k_l}{k_l+l\gamma} \frac{l(b+1)\gamma}{bk_{l-1}} \right) \right) \right). \quad (\text{S2.28})$$

S3. OPTIMAL FEEDBACK WHEN PROTEIN DOES NOT DEGRADE

As mentioned in the main text, our objective is to minimize $\langle FPT^2 \rangle$ such that $\langle FPT \rangle$ is fixed. For calculation purposes, we will denote this constraint as $\langle FPT \rangle = t_{opt}$. Let m represents the Lagrange's multiplier, then we define the following objective function

$$\phi := \langle FPT^2 \rangle + m (\langle FPT \rangle - t_{opt}). \quad (\text{S3.1})$$

The optimization problem is solved in two steps. First, we determine the critical points. Second, we find the critical point corresponding to a global minimum.

Determining the critical points requires the following system of equations to be solved

$$\frac{\partial \langle FPT^2 \rangle}{\partial k_i} = m \frac{\partial \langle FPT \rangle}{\partial k_i}, 0 \leq i \leq X-1, \quad (\text{S3.2})$$

$$\langle FPT \rangle = t_{opt}. \quad (\text{S3.3})$$

The expressions of the first-two moments of FPT when protein does not degrade are given

by equation (16) in the main text. These are

$$\langle FPT \rangle = \frac{1}{k_0} + \frac{1}{b} \sum_{i=0}^{X-1} \frac{1}{k_i}, \quad (\text{S3.4})$$

$$\langle FPT^2 \rangle = \frac{2}{b^2} \left(\frac{\tau_0}{bk_0} + \sum_{i=0}^{X-1} \frac{\tau_i}{k_i} \right), \text{ where} \quad (\text{S3.5})$$

$$\tau_i := \frac{b}{k_i} + \sum_{j=i}^{X-1} \frac{1}{k_j}. \quad (\text{S3.6})$$

The optimization problem in equations (S3.2),(S3.3) requires calculation of the first order derivatives of $\langle FPT \rangle$ and $\langle FPT^2 \rangle$. The derivatives of $\langle FPT \rangle$ with respect to k_i 's are given by

$$\frac{\partial \langle FPT \rangle}{\partial k_0} = \frac{b+1}{b} \left(-\frac{1}{k_0^2} \right); \quad \frac{\partial \langle FPT \rangle}{\partial k_i} = \frac{1}{b} \left(-\frac{1}{k_i^2} \right), \quad 1 \leq i \leq X-1. \quad (\text{S3.7})$$

Similarly, the derivative of $\langle FPT^2 \rangle$ are

$$\frac{\partial \langle FPT^2 \rangle}{\partial k_0} = 2 \left(\frac{1}{b} \right)^2 \left(-\frac{2(b+1)^2}{k_0^3} - \frac{b+1}{k_0^2} \sum_{j=1}^{X-1} \frac{1}{k_j} \right), \quad (\text{S3.8})$$

$$\frac{\partial \langle FPT^2 \rangle}{\partial k_i} = 2 \left(\frac{1}{b} \right)^2 \left(\frac{b+1}{k_0} + \frac{2(b+1)}{k_i} + \sum_{\substack{j=1 \\ j \neq i}}^{X-1} \frac{1}{k_j} \right) \left(\frac{-1}{k_i^2} \right), \quad 1 \leq i \leq X-1. \quad (\text{S3.9})$$

Therefore, the system of equations to be solved becomes

$$2 \left(\frac{1}{b} \right)^2 \left(-\frac{2(b+1)^2}{k_0^3} - \frac{b+1}{k_0^2} \sum_{j=1}^{X-1} \frac{1}{k_j} \right) = m \frac{b+1}{b} \left(-\frac{1}{k_0^2} \right), \quad (\text{S3.10})$$

$$2 \left(\frac{1}{b} \right)^2 \left(\frac{b+1}{k_0} + \frac{2(b+1)}{k_i} + \sum_{\substack{j=1 \\ j \neq i}}^{X-1} \frac{1}{k_j} \right) \left(\frac{-1}{k_i^2} \right) = m \frac{1}{b} \left(-\frac{1}{k_i^2} \right), \quad 1 \leq i \leq X-1, \quad (\text{S3.11})$$

$$\frac{1}{b} \left(\frac{b+1}{k_0} + \sum_{j=1}^{X-1} \frac{1}{k_j} \right) = t_{opt}. \quad (\text{S3.12})$$

Furthermore, we want the solution such that none of the transcription rate is zero. This

simplifies the system of equations to

$$\frac{2}{b} \left(\frac{2(b+1)}{k_0} + \sum_{j=1}^{X-1} \frac{1}{k_j} \right) = m, \quad (\text{S3.13})$$

$$\frac{2}{b} \left(\frac{b+1}{k_0} + \frac{2(b+1)}{k_i} + \sum_{\substack{j=1 \\ j \neq i}}^{X-1} \frac{1}{k_j} \right) = m, \quad 1 \leq i \leq X-1, \quad (\text{S3.14})$$

$$\frac{1}{b} \left(\frac{b+1}{k_0} + \sum_{j=1}^{X-1} \frac{1}{k_j} \right) = t_{opt}. \quad (\text{S3.15})$$

Solution to these equations is given by

$$k_0 = \frac{(b+1)(2b+X)}{(2b+1)bt_{opt}} \quad (\text{S3.16})$$

$$k_i = \frac{2b+1}{b+1} k_0 = \frac{2b+X}{bt_{opt}}, \quad 1 \leq i \leq X-1, \quad (\text{S3.17})$$

$$m = \frac{2(b+1)(4b+X+1)}{(2b+1)bk_0} = \frac{2t_{opt}(4b+X+1)}{2b+X}. \quad (\text{S3.18})$$

We have calculated the critical point for the optimization problem. However, it needs to be checked whether its an minimum or maximum. For this purpose, we consider the bordered Hessian as follows.

$$D_\phi = \begin{bmatrix} \frac{\partial^2 \phi}{\partial m^2} & \frac{\partial^2 \phi}{\partial m \partial k_0} & \frac{\partial^2 \phi}{\partial m \partial k_1} & \cdots & \frac{\partial^2 \phi}{\partial m \partial k_{X-1}} \\ \frac{\partial^2 \phi}{\partial k_0 \partial m} & \frac{\partial^2 \phi}{\partial k_0^2} & \frac{\partial^2 \phi}{\partial k_0 \partial k_1} & \cdots & \frac{\partial^2 \phi}{\partial k_0 \partial k_{X-1}} \\ \frac{\partial^2 \phi}{\partial k_1 \partial m} & \frac{\partial^2 \phi}{\partial k_1 \partial k_0} & \frac{\partial^2 \phi}{\partial k_1^2} & \cdots & \frac{\partial^2 \phi}{\partial k_1 \partial k_{X-1}} \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ \frac{\partial^2 \phi}{\partial k_{X-1} \partial m} & \frac{\partial^2 \phi}{\partial k_{X-1} \partial k_0} & \frac{\partial^2 \phi}{\partial k_{X-1} \partial k_1} & \cdots & \frac{\partial^2 \phi}{\partial k_{X-1}^2} \end{bmatrix}. \quad (\text{S3.19})$$

We will show that all the principal minors of this matrix are negative. To start with, let us first determine the second order derivatives of $\langle FPT \rangle$.

$$\frac{\partial^2 \langle FPT \rangle}{\partial k_0^2} = \frac{\partial}{\partial k_0} \left(-\frac{b+1}{bk_0^2} \right) = \frac{2(b+1)}{bk_0^3}, \quad (\text{S3.20})$$

$$\frac{\partial^2 \langle FPT \rangle}{\partial k_0 \partial k_i} = \frac{\partial}{\partial k_0} \left(-\frac{1}{bk_i^2} \right) = 0, \quad 1 \leq i \leq X-1, \quad (\text{S3.21})$$

$$\frac{\partial^2 \langle FPT \rangle}{\partial k_i \partial k_0} = \frac{\partial}{\partial k_i} \left(-\frac{b+1}{bk_0^2} \right) = 0, \quad 1 \leq i \leq X-1, \quad (\text{S3.22})$$

$$\frac{\partial^2 \langle FPT \rangle}{\partial k_i \partial k_j} = \frac{\partial}{\partial k_i} \left(-\frac{1}{bk_j^2} \right) = 0, \quad 1 \leq i, j \leq X-1, i \neq j \quad (\text{S3.23})$$

$$\frac{\partial^2 \langle FPT \rangle}{\partial k_i^2} = \frac{\partial}{\partial k_i} \left(-\frac{1}{bk_i^2} \right) = \frac{2}{bk_i^3}, \quad 1 \leq i \leq X-1, \quad (\text{S3.24})$$

Similarly, the derivatives for $\langle FPT^2 \rangle$ are given by

$$\begin{aligned} \frac{\partial^2 \langle FPT^2 \rangle}{\partial k_0^2} &= \frac{\partial}{\partial k_0} \left(2 \left(\frac{1}{b} \right)^2 \left(-\frac{2(b+1)^2}{k_0^3} - \frac{b+1}{k_0^2} \sum_{j=1}^{X-1} \frac{1}{k_j} \right) \right) \\ &= 2 \left(\frac{1}{b} \right)^2 \left(\frac{6(b+1)^2}{k_0^4} + \frac{2(b+1)}{k_0^3} \sum_{j=1}^{X-1} \frac{1}{k_j} \right), \end{aligned} \quad (\text{S3.25})$$

$$\begin{aligned} \frac{\partial^2 \langle FPT^2 \rangle}{\partial k_0 \partial k_i} &= \frac{\partial}{\partial k_0} \left(2 \left(\frac{1}{b} \right)^2 \left(\frac{b+1}{k_0} + \frac{2(b+1)}{k_i} + \sum_{\substack{j=1 \\ j \neq i}}^{X-1} \frac{1}{k_j} \right) \left(\frac{-1}{k_i^2} \right) \right) \\ &= 2 \left(\frac{1}{b} \right)^2 \frac{b+1}{k_0^2 k_i^2}, 1 \leq i \leq X-1, \end{aligned} \quad (\text{S3.26})$$

$$\begin{aligned} \frac{\partial^2 \langle FPT^2 \rangle}{\partial k_i \partial k_0} &= \frac{\partial}{\partial k_i} \left(2 \left(\frac{1}{b} \right)^2 \left(-\frac{2(b+1)^2}{k_0^3} - \frac{b+1}{k_0^2} \sum_{j=1}^{X-1} \frac{1}{k_j} \right) \right) \\ &= 2 \left(\frac{1}{b} \right)^2 \frac{b+1}{k_0^2 k_i^2}, 1 \leq i \leq X-1, \end{aligned} \quad (\text{S3.27})$$

$$\begin{aligned} \frac{\partial^2 \langle FPT^2 \rangle}{\partial k_i \partial k_j} &= \frac{\partial}{\partial k_i} \left(2 \left(\frac{1}{b} \right)^2 \left(\frac{b+1}{k_0} + \frac{2(b+1)}{k_j} + \sum_{\substack{l=1 \\ l \neq j}}^{X-1} \frac{1}{k_l} \right) \left(\frac{-1}{k_j^2} \right) \right) \\ &= 2 \left(\frac{1}{b} \right)^2 \frac{1}{k_i^2 k_j^2}, 1 \leq i, j \leq X-1, i \neq j, \end{aligned} \quad (\text{S3.28})$$

$$\begin{aligned} \frac{\partial^2 \langle FPT^2 \rangle}{\partial k_i^2} &= \frac{\partial}{\partial k_i} \left(2 \left(\frac{1}{b} \right)^2 \left(\frac{b+1}{k_0} + \frac{2(b+1)}{k_i} + \sum_{\substack{j=1 \\ j \neq i}}^{X-1} \frac{1}{k_j} \right) \left(\frac{-1}{k_i^2} \right) \right) \\ &= 2 \left(\frac{1}{b} \right)^2 \left(\frac{2(b+1)}{k_0 k_i^3} + \frac{6(b+1)}{k_i^4} + \frac{2}{k_i^3} \sum_{j=1}^{X-1} \frac{1}{k_j} \right), 1 \leq i \leq X-1. \end{aligned} \quad (\text{S3.29})$$

We can now determine the elements of the bordered Hessian matrix in equation (S3.19) computed at the solution given by equations (S3.16)-(S3.17).

$$\frac{\partial^2 \phi}{\partial m^2} = \frac{\partial^2 (\langle FPT^2 \rangle + m (\langle FPT \rangle - t_{opt}))}{\partial m^2} = 0, \quad (\text{S3.30})$$

$$\frac{\partial^2 \phi}{\partial m \partial k_0} = \frac{\partial^2 \phi}{\partial k_0 \partial m} = \frac{\partial \langle FPT \rangle}{\partial k_0} = \frac{b+1}{b} \left(-\frac{1}{k_0^2} \right). \quad (\text{S3.31})$$

$$\frac{\partial^2 \phi}{\partial m \partial k_i} = \frac{\partial^2 \phi}{\partial k_i \partial m} = \frac{\partial \langle FPT \rangle}{\partial k_i} = \frac{1}{b} \left(-\frac{1}{k_i^2} \right). \quad (\text{S3.32})$$

$$\frac{\partial^2 \phi}{\partial k_0^2} = \frac{\partial^2 (\langle FPT^2 \rangle + m (\langle FPT \rangle - t_{opt}))}{\partial k_0^2} \quad (\text{S3.33})$$

$$= 2 \left(\frac{1}{b} \right)^2 \left(\frac{6(b+1)^2}{k_0^4} + \frac{2(b+1)}{k_0^3} \sum_{j=1}^{X-1} \frac{1}{k_j} \right) + m \frac{2(b+1)}{bk_0^3} \quad (\text{S3.34})$$

$$= \frac{4(b+1)^2(10b+2X+3)}{(2b+1)b^2k_0^4} \quad (\text{S3.35})$$

$$\frac{\partial^2 \phi}{\partial k_i^2} = \frac{\partial^2 (\langle FPT^2 \rangle + m (\langle FPT \rangle - t_{opt}))}{\partial k_i^2} \quad (\text{S3.36})$$

$$= 2 \left(\frac{1}{b} \right)^2 \left(\frac{2(b+1)}{k_0 k_i^3} + \frac{6(b+1)}{k_i^4} + \frac{2}{k_i^3} \sum_{j=1}^{X-1} \frac{1}{k_j} \right) + m \frac{2}{bk_i^3} \quad (\text{S3.37})$$

$$= \frac{4(b+1)^4(9b+2X+4)}{(2b+1)^4 b^2 k_0^4}. \quad (\text{S3.38})$$

$$\frac{\partial^2 \phi}{\partial k_0 k_i} = \frac{\partial^2 \phi}{\partial k_i k_0} = \frac{\partial^2 (\langle FPT^2 \rangle + m (\langle FPT \rangle - t_{opt}))}{\partial k_0 \partial k_i} \quad (\text{S3.39})$$

$$= 2 \left(\frac{1}{b} \right)^2 \frac{b+1}{k_0^2 k_i^2} \quad (\text{S3.40})$$

$$= \frac{2(b+1)^3}{(2b+1)^2 b^2 k_0^4} \quad (\text{S3.41})$$

$$\frac{\partial^2 \phi}{\partial k_j k_i} = \frac{\partial^2 \phi}{\partial k_i k_j} = \frac{\partial^2 (\langle FPT^2 \rangle + m (\langle FPT \rangle - t_{opt}))}{\partial k_j \partial k_i} \quad (\text{S3.42})$$

$$= 2 \left(\frac{1}{b} \right)^2 \frac{1}{k_j^2 k_i^2} \quad (\text{S3.43})$$

$$= \frac{2(b+1)^2}{(2b+1)^2 b^2 k_0^4} \quad (\text{S3.44})$$

It can be noted elements of D_ϕ are from a set six quantities. Defining

$$\begin{aligned} q_1 &:= \frac{b+1}{b} \left(-\frac{1}{k_0^2} \right), & q_2 &:= \frac{1}{b} \left(-\frac{1}{k_0^2} \right), \\ q_3 &:= \frac{4(b+1)^2(10b+2X+3)}{(2b+1)b^2k_0^4}, & q_4 &:= \frac{2(b+1)^3}{(2b+1)^2b^2k_0^4}, \\ q_5 &:= \frac{4(b+1)^4(9b+2X+4)}{(2b+1)^4b^2k_0^4}, & q_6 &:= \frac{2(b+1)^2}{(2b+1)^2b^2k_0^4} \end{aligned} \quad (\text{S3.45})$$

we can write D_ϕ as

$$D_\phi = \begin{bmatrix} 0 & q_1 & q_2 & \cdots & q_2 \\ q_1 & q_3 & q_4 & \cdots & q_4 \\ q_2 & q_4 & q_5 & \cdots & q_6 \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ q_2 & q_4 & q_6 & \cdots & q_5 \end{bmatrix}. \quad (\text{S3.46})$$

Let us denote by $K(n)$ the principal minor of the matrix D_ϕ of size $n \times n$. It can be easily seen that $K(1) = 0$, $K(2) = 0 - q_1^2 < 0$ and $K(3) = -q_2^2q_3 + 2q_1q_2q_4 - q_1^2q_5$. For $4 \leq n \leq X$, we perform the following two elementary operations on D_ϕ

- $col_r = col_r - col_{r-1}$
- $row_r = row_i - row_{r-1}$

for $r = n, n-1, \dots, 1$. This yields

$$K(n) = 2(q_5 - q_6)K(n-1) - (q_5 - q_6)^2K(n-2), \quad 4 \leq n \leq X. \quad (\text{S3.47})$$

The solution to the above recursive equation is given by

$$K(n) = -(q_5 - q_6)^{n-3} \left((n-2) (q_2^2q_3 - 2q_1q_2q_4 + q_1^2q_6) + (q_1^2q_6 + q_1^2q_5) \right). \quad (\text{S3.48})$$

It can be easily checked that $K(n)$ is negative because $q_5 > q_6$, $q_2^2q_3 - 2q_1q_2q_4 + q_1^2q_6 > 0$ and $q_1^2q_6 + q_1^2q_5 > 0$. This proves that the critical point indeed corresponds to a minimum.

S4. OPTIMAL FEEDBACK STRATEGY IN PRESENCE OF PROTEIN DEGRADATION

In previous sections, we have derived analytical expression of the optimal feedback strategy that minimizes $\langle FPT^2 \rangle$ such that $\langle FPT \rangle$ is constant. As the expressions of $\langle FPT \rangle$

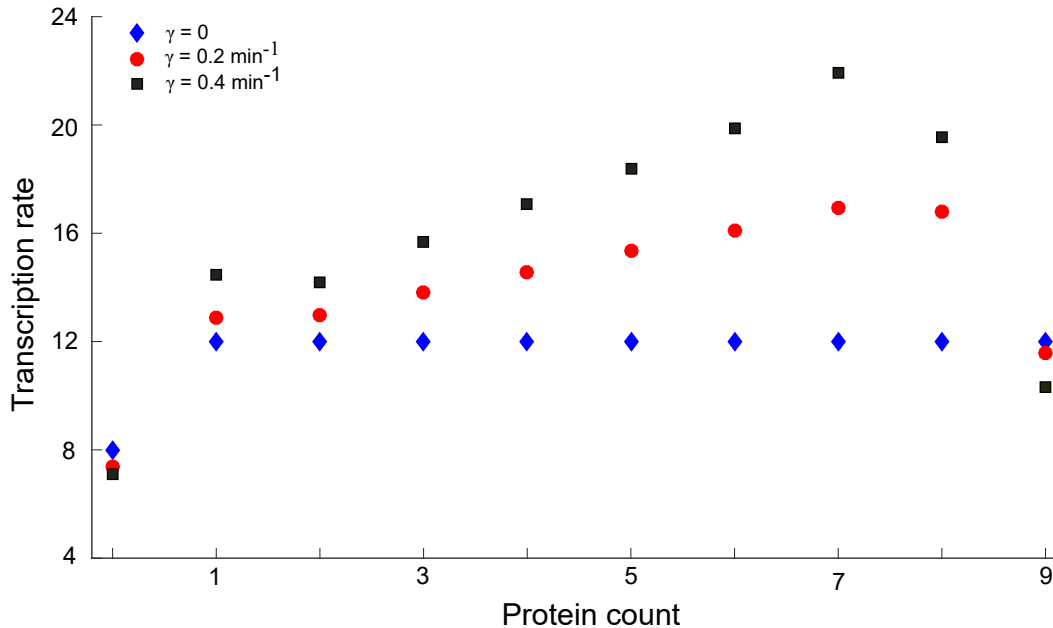


FIG. S4.1. *Optimal feedback strategy for unstable protein.* The optimal transcription rates obtained via numerical optimization for different values of degradation rate are shown. The event threshold is assumed to be 10 molecules, and the mean FPT is constrained to be 1 minute.

and $\langle FPT^2 \rangle$ in equations (S2.20) and (S2.28) are too convoluted to solve for optimal transcription rates analytically, we take a numerical approach. For this purpose, we fixed the threshold $X = 10$ molecules, and mean burst size $b = 1$ molecules. Using numerical solvers, we searched the parameter space of the transcription rates $k_i, i \in \{0, 1, \dots, 9\}$ for various values of the protein degradation rate γ . As shown in Fig. S4.1, when $\gamma = 0$, the transcription rates are equal except for the first one. This is consistent with the expressions in (S3.16)-(S3.17). Further, as the protein degradation rate is increased, the transcription rates first increase and then decrease, suggesting a mixed feedback strategy.

To keep the results biologically meaningful, we assume that the feedbacks are implemented as Hill functions. One simplest implementation of this would be a linear form of the feedback

$$k_i = c_1 + c_2 i, \quad (\text{S4.1})$$

as considered in the analysis in the main text. Here, we show the results for having a nonlinear form of Hill functions wherein a positive feedback is implemented as below

$$k_i = k_{max} \left(r + (1 - r) \frac{(ci)^H}{1 + (ci)^H} \right), \quad (\text{S4.2})$$

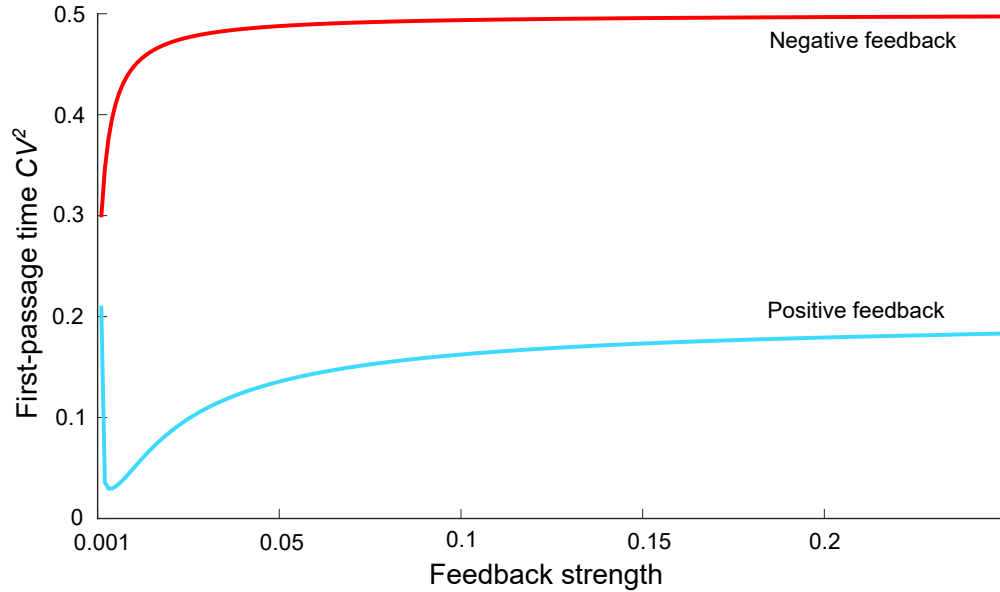


FIG. S4.2. *Optimal feedback strategy in form of a Hill function.* The transcription rates are assumed to follow the Hill function forms as given by (S4.2) (positive feedback) and (S4.3) (negative feedback). For a negative feedback, CV_{FPT}^2 increases as the feedback strength increased. For a positive feedback, CV_{FPT}^2 first decreases as feedback strength increases, hits a minimum at a certain feedback strength, and increases after that. The parameters used are $b = 1$ molecule, $r = 0.05$, and $H = 1$.

where k_{max} represents the maximum possible transcription rate, $0 < r < 1$ is a constant corresponding to the minimum transcription rate ($k_{max}r$), c is the feedback strength (note that $k_i = k_{max}/2$ when $i=1/c$), and H denotes the Hill coefficient. Similarly, a negative feedback is implemented via

$$k_i = k_{max} \left(r + \frac{1 - r}{1 + (ci)^H} \right). \quad (\text{S4.3})$$

The results of these Hill function implementations are same as shown in the main text for the linear form of transcription rates. We show one example of this in Fig. S4.2. One can see that when the protein degradation is allowed, the negative feedback leads to increase in noise as its feedback strength is increased. For the positive feedback, the noise hits a minimum for a certain feedback strength. Here the mean FPT is kept constant by appropriately tuning the parameter k_{max} .

One can see that there are more parameters in these forms of feedback, and various parameters (or their combinations) can be tuned to keep the mean fixed for a given feedback

strength, the resulting analysis is complex. This is the rationale behind using the simplest possible form (linear) of the feedback which gives useful insight.

S5. OPTIMAL FEEDBACK STRENGTH FOR LINEAR FORM OF FEEDBACK IN PRESENCE OF PROTEIN DEGRADATION

In this section, we explore the optimal feedback strength c_2 as the protein degradation rate is varied. As in the main text, the feedbacks are assumed to follow a linear form given by $k_i = c_1 \pm c_2 i$.

When the protein does not decay, as expected from theory, we get the optimal $c_2 = 0$. As protein decay is considered, a positive feedback acts as the optimal feedback strategy. The optimal c_2 multiplied by the mean burst size b takes values close to the degradation rate γ as shown in Fig. S5.1 (left). Interestingly, the protein trajectories generated by the optimal (positive) feedback in the case of protein degradation mimic those generated by the optimal feedback (no feedback) when protein did not decay. For their average dynamics to follow each other, a deterministic analysis reveals that as $\dot{x} = b(c_1 + c_2 x) - \gamma x$, bc_2 should be equal to γ . The slight difference between them appears to be due to the fact that a stochastic mean and a deterministic mean usually differ from each other.

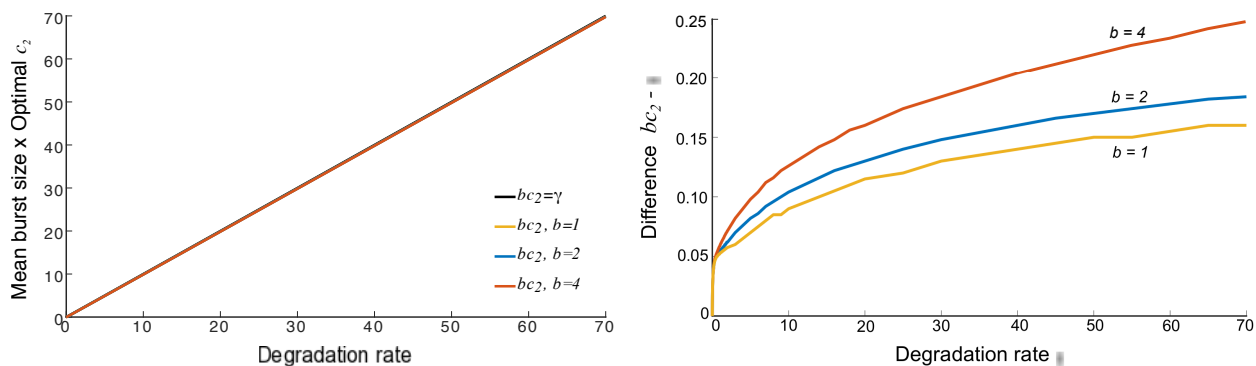


FIG. S5.1. **Difference between the degradation rate γ and optimal feedback strength c_2 multiplied by the mean burst size b :** *Left:* The quantities are plotted for a range of degradation rates. It can be seen that they remain close to each other for increase in the degradation rate, and for several burst sizes. *Right:* The difference between the quantities is plotted as γ is changed. For a given mean burst size b , the difference $bc_2 - \gamma$ increases with increase in γ . For a given degradation rate, this difference also increases as the mean burst size b is increased.

We also explored the difference $bc_2 - \gamma$ as γ varies. Intriguingly, if the mean burst size b is kept constant, we observe that the difference increases for increase in γ . Further, for a given degradation rate, increasing b leads to increase in the difference $bc_2 - \gamma$. The reason for this is not clear to us yet, though it appears that in this case the stochastic description of the dynamics shows more deviations from the deterministic dynamics.

S6. OPTIMAL FEEDBACK FOR STABLE PROTEIN IN PRESENCE OF EXTRINSIC NOISE

In this section, assuming that the protein does not degrade, we investigate how the optimal regulation strategy deviates from a no feedback in presence of a static extrinsic noise. We consider two possibilities here: one, the extrinsic noise affects the mean burst size; two, the extrinsic noise affects the transcription rate. For the first case, we assume that the mean burst size is drawn from an arbitrary positive-valued distribution. The second case is analyzed by assuming that a factor Z multiplies with the transcription rates resulting in an effective transcription rate when $x(t) = i$ to be $k_i Z$.

S6-a. Optimal feedback when the mean burst size is drawn from a distribution

Let us assume the mean burst size a random variable with probability density function $f_b(\beta)$. Thus, the number of proteins in a burst are geometrically distributed with mean b where b itself is a random variable. The mean FPT can be computed as

$$\langle FPT | b = \beta \rangle = \frac{1}{\beta} \left(\frac{\beta + 1}{k_0} + \sum_{i=1}^{X-1} \frac{1}{k_i} \right), \quad (\text{S6.1})$$

$$\implies \langle FPT \rangle = \left(1 + \left\langle \frac{1}{b} \right\rangle \right) \frac{1}{k_0} + \left\langle \frac{1}{b} \right\rangle \sum_{i=1}^{X-1} \frac{1}{k_i}. \quad (\text{S6.2})$$

Similarly, the second order moment of FPT is given by

$$\langle FPT^2 | b = \beta \rangle = \left(\frac{1}{\beta} \right)^2 \left(\frac{\beta + 1}{k_0} \left(\frac{\beta + 1}{k_0} + \sum_{i=1}^{X-1} \frac{1}{k_i} \right) + \sum_{i=1}^{X-1} \frac{1}{k_i} \left(\frac{\beta + 1}{k_i} + \sum_{j=i+1}^{X-1} \frac{1}{k_j} \right) \right), \quad (\text{S6.3})$$

$$\begin{aligned} \Rightarrow \langle FPT^2 \rangle &= \left(\left\langle \frac{1}{b^2} \right\rangle + 2 \left\langle \frac{1}{b} \right\rangle + 1 \right) \frac{1}{k_0^2} + \left(\left\langle \frac{1}{b^2} \right\rangle + \left\langle \frac{1}{b} \right\rangle \right) \left(\frac{1}{k_0} \sum_{i=1}^{X-1} \frac{1}{k_i} + \sum_{i=1}^{X-1} \frac{1}{k_i^2} \right) \\ &\quad + \left\langle \frac{1}{b^2} \right\rangle \sum_{i=1}^{X-1} \frac{1}{k_i} \sum_{j=i+1}^{X-1} \frac{1}{k_j}. \end{aligned} \quad (S6.4)$$

Defining $\alpha_1 = \langle \frac{1}{b} \rangle$ and $\alpha_2 = \langle \frac{1}{b^2} \rangle$, we have

$$\langle FPT \rangle = (1 + \alpha_1) \frac{1}{k_0} + \alpha_1 \sum_{i=1}^{X-1} \frac{1}{k_i}, \quad (S6.5)$$

$$\langle FPT^2 \rangle = (1 + 2\alpha_1 + \alpha_2) \frac{1}{k_0^2} + (\alpha_1 + \alpha_2) \left(\frac{1}{k_0} \sum_{i=1}^{X-1} \frac{1}{k_i} + \sum_{i=1}^{X-1} \frac{1}{k_i^2} \right) + \alpha_2 \sum_{i=1}^{X-1} \frac{1}{k_i} \sum_{j=i+1}^{X-1} \frac{1}{k_j} \quad (S6.6)$$

The derivatives of $\langle FPT \rangle$ with respect to k_i 's are given by

$$\frac{\partial \langle FPT \rangle}{\partial k_0} = (1 + \alpha_1) \left(-\frac{1}{k_0^2} \right); \quad \frac{\partial \langle FPT \rangle}{\partial k_i} = \alpha_1 \left(-\frac{1}{k_i^2} \right), \quad 1 \leq i \leq X - 1. \quad (S6.7)$$

Similarly, the derivative of $\langle FPT^2 \rangle$ are

$$\frac{\partial \langle FPT^2 \rangle}{\partial k_0} = 2 \left(-\frac{2(1 + 2\alpha_1 + \alpha_2)}{k_0^3} - \frac{(\alpha_1 + \alpha_2)}{k_0^2} \sum_{j=1}^{X-1} \frac{1}{k_j} \right); \quad (S6.8)$$

$$\frac{\partial \langle FPT^2 \rangle}{\partial k_i} = 2 \left(\frac{\alpha_1 + \alpha_2}{k_0} + \frac{2(\alpha_1 + \alpha_2)}{k_i} + \alpha_2 \sum_{\substack{j=1 \\ j \neq i}}^{X-1} \frac{1}{k_j} \right) \left(\frac{-1}{k_i^2} \right), \quad 1 \leq i \leq X - 1. \quad (S6.9)$$

To find a critical point, we have to solve the following system of equations

$$2 \left(-\frac{2(1 + 2\alpha_1 + \alpha_2)}{k_0^3} - \frac{(\alpha_1 + \alpha_2)}{k_0^2} \sum_{j=1}^{X-1} \frac{1}{k_j} \right) = m(1 + \alpha_1) \left(-\frac{1}{k_0^2} \right) \quad (S6.10)$$

$$2 \left(\frac{\alpha_1 + \alpha_2}{k_0} + \frac{2(\alpha_1 + \alpha_2)}{k_i} + \alpha_2 \sum_{\substack{j=1 \\ j \neq i}}^{X-1} \frac{1}{k_j} \right) \left(\frac{-1}{k_i^2} \right) = m\alpha_1 \left(-\frac{1}{k_i^2} \right), \quad 1 \leq i \leq X - 1, \quad (S6.11)$$

$$(1 + \alpha_1) \frac{1}{k_0} + \alpha_1 \sum_{i=1}^{X-1} \frac{1}{k_i} = t_{opt}. \quad (S6.12)$$

Assuming that $k_0, k_1, \dots \neq 0$, we get

$$2 \left(\frac{2(1 + 2\alpha_1 + \alpha_2)}{k_0} + (\alpha_1 + \alpha_2) \sum_{j=1}^{X-1} \frac{1}{k_j} \right) = m(1 + \alpha_1) \quad (\text{S6.13})$$

$$2 \left(\frac{\alpha_1 + \alpha_2}{k_0} + \frac{2(\alpha_1 + \alpha_2)}{k_i} + \alpha_2 \sum_{\substack{j=1 \\ j \neq i}}^{X-1} \frac{1}{k_j} \right) = m\alpha_1, 1 \leq i \leq X-1, \quad (\text{S6.14})$$

$$(1 + \alpha_1) \frac{1}{k_0} + \alpha_1 \sum_{i=1}^{X-1} \frac{1}{k_i} = t_{opt}. \quad (\text{S6.15})$$

Solution to above system of equations gives

$$k_i = k = k_0 \zeta, 1 \leq i \leq X-1, \quad (\text{S6.16})$$

$$k_0 = \frac{X\alpha_1((1 + \alpha_1)^2 + 2\alpha_2) - 3\alpha_1^3 + \alpha_2 + 2\alpha_1^2 + \alpha_1^2\alpha_2 + \alpha_1}{(X(\alpha_2 - \alpha_1^2) + \alpha_1(2 + 3\alpha_1 + \alpha_2))t_{opt}}, \quad (\text{S6.17})$$

$$\zeta = \frac{X(\alpha_2 - \alpha_1^2) + 3\alpha_1^2 + 2\alpha_1 + \alpha_1\alpha_2}{3\alpha_1^2 + \alpha_1 + \alpha_1\alpha_2 - \alpha_2}. \quad (\text{S6.18})$$

These equations reduce to our previous results of having a constant mean burst size b when $\alpha_1 = 1/b$, $\alpha_2 = 1/b^2$ are used.

S6-b. Optimal regulation when extrinsic factor affects the transcription rate

We consider an extrinsic factor Z with a positive-valued arbitrary distribution $f_Z(z)$. This factor is assumed to be static, i.e., it does not vary over the time scale of the event of interest. Further we assume that it affects the transcription rates in a multiplicative fashion. The first-passage time mean in this case can be written as

$$\langle FPT | Z = z \rangle = \frac{1}{z} \left(\frac{1}{k_0} + \frac{1}{b} \sum_{i=0}^{X-1} \frac{1}{k_i} \right) \implies \langle FPT \rangle = \left\langle \frac{1}{z} \right\rangle \left(\frac{1}{k_0} + \frac{1}{b} \sum_{i=0}^{X-1} \frac{1}{k_i} \right). \quad (\text{S6.19})$$

Likewise the second order moment can be written as

$$\langle FPT^2 \rangle = \left\langle \frac{1}{z^2} \right\rangle \frac{2}{b^2} \left(\frac{\tau_0}{bk_0} + \sum_{i=0}^{X-1} \frac{\tau_i}{k_i} \right), \tau_i := \frac{b}{k_i} + \sum_{j=i}^{X-1} \frac{1}{k_j}. \quad (\text{S6.20})$$

Solving the constrained optimization problem of minimizing $\langle FPT^2 \rangle$ constrained to $\langle FPT \rangle = t_{opt}$ in this case simplifies to solving the following system of equations

$$\frac{\langle 1/z^2 \rangle}{\langle 1/z \rangle} \frac{2}{b} \left(\frac{2(b+1)}{k_0} + \sum_{j=1}^{X-1} \frac{1}{k_j} \right) = m, \quad (\text{S6.21})$$

$$\frac{\langle 1/z^2 \rangle}{\langle 1/z \rangle} \frac{2}{b} \left(\frac{b+1}{k_0} + \frac{2(b+1)}{k_i} + \sum_{\substack{j=1 \\ j \neq i}}^{X-1} \frac{1}{k_j} \right) = m, \quad 1 \leq i \leq X-1, \quad (\text{S6.22})$$

$$\left\langle \frac{1}{z} \right\rangle \frac{1}{b} \left(\frac{b+1}{k_0} + \sum_{j=1}^{X-1} \frac{1}{k_j} \right) = t_{opt}, \quad (\text{S6.23})$$

where m represents the Lagrange's multiplier. Solution to these equations gives

$$k_0 = \frac{(b+1)(2b+X) \langle \frac{1}{z} \rangle}{(2b+1)bt_{opt}} \quad (\text{S6.24})$$

$$k_i = \frac{2b+1}{b+1} k_0 = \frac{(2b+X) \langle \frac{1}{z} \rangle}{bt_{opt}}, \quad 1 \leq i \leq X-1. \quad (\text{S6.25})$$

S7. OPTIMAL FEEDBACK STRATEGY WHEN BURST SIZE IS DRAWN FROM A POISSON DISTRIBUTION

In the main paper, the mRNA degradation process is assumed to be memoryless (exponential), and consequently the burst of proteins is assumed to follow a geometric distribution. However, in the limit when the mRNA degradation process is deterministic, the burst size distribution becomes Poisson. For this reason, we investigate how the optimal feedback strategy changes in the case when the burst follows a Poisson distribution.

The production and degradation of the protein (similar to equation (2) in main text) is governed by following probabilities

$$\mathbb{P}(x(t+dt) = i+B | x(t) = i) = k_i dt, \quad (\text{S7.1})$$

$$\mathbb{P}(x(t+dt) = i-1 | x(t) = i) = i\gamma dt, \quad (\text{S7.2})$$

where the burst size B follows a Poisson distribution given by

$$\mathbb{P}(B = j) = \frac{b^j}{j!} e^{-b}. \quad (\text{S7.3})$$

Here b , as before, represents the mean burst size, i.e., the average number of protein molecules produced by one mRNA.

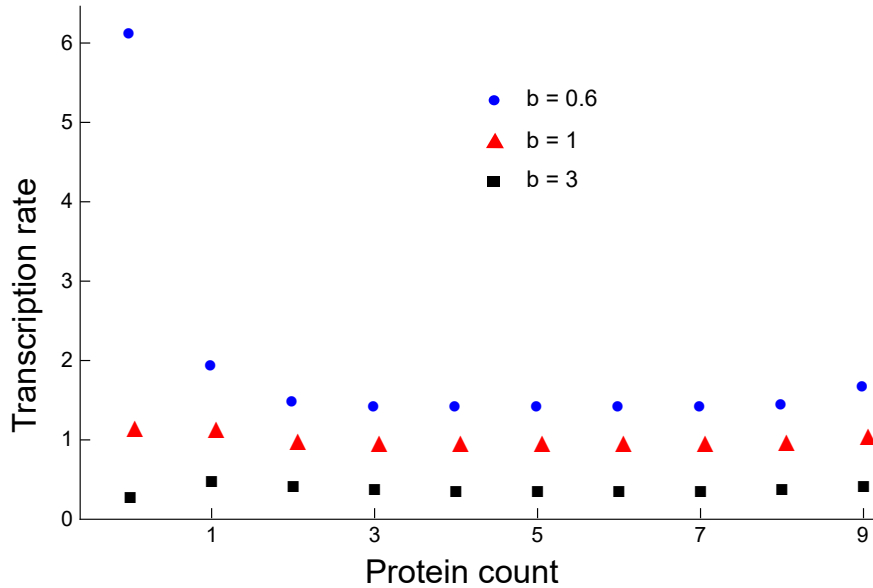


FIG. S7.1. *Optimal feedback strategy for Poisson distributed burst size.* The optimal transcription rates obtained via numerical optimization for different values of mean burst size are shown. The event threshold is assumed to be 10 molecules, and the mean *FPT* is constrained to be 10 minutes.

One can carry out the first-passage time calculations in the same manner as done for the geometric burst size case. It turns out that the first two moments can be compactly written as

$$\langle FPT \rangle = \sum_{i=0}^{X-1} \frac{\alpha_i}{k_i}, \quad (S7.4)$$

$$\langle FPT^2 \rangle = 2 \sum_{i=0}^{X-1} \frac{\tau_i}{k_i}, \quad (S7.5)$$

where

$$\tau_i = \sum_{j=i}^{X-1} \frac{\alpha_{j-i}}{k_j}. \quad (S7.6)$$

The coefficients α_i , $i \in \{0, 1, 2, \dots, X-1\}$ are defined as

$$\alpha_0 = 1, \quad \alpha_i = \frac{e^b}{e^b - 1} \frac{b^i}{i! (e^b - 1)^i} \sum_{m=0}^{i-1} a[i, m] e^{mb}, \quad (S7.7)$$

with $a[r, s]$ represents an Eulerian number whose expression is

$$a[r, s] = \sum_{i=0}^{s+1} (-1)^i \binom{r+1}{i} (s+1-i)^r. \quad (S7.8)$$

We performed numerical optimization with respect to parameters k_i 's for threshold $X = 10$ to see the form of the optimal feedback strategy such that $\langle FPT^2 \rangle$ is minimized, with constraint $\langle FPT \rangle = t_{opt} = 10$ minutes. The results show that while the optimal transcription rates are not equal (i.e., no feedback mechanism in strict sense), they are fairly close to each other for mean burst sizes of 1 and 3 molecules. For mean burst size of 0.6 molecules, the first transcription rate when protein count is zero comes out to be significantly higher than other transcription rates which are more or less close to each other (Fig. S7.1). These results suggest that while the optimal feedback strategy deviates from a no feedback strategy with the underlying distribution of the burst size, it appears to remain close to a no feedback strategy.

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